



# Backtracking and Branch and Bound

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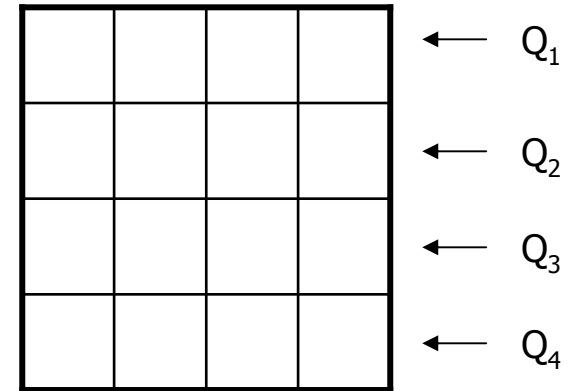
# Backtracking

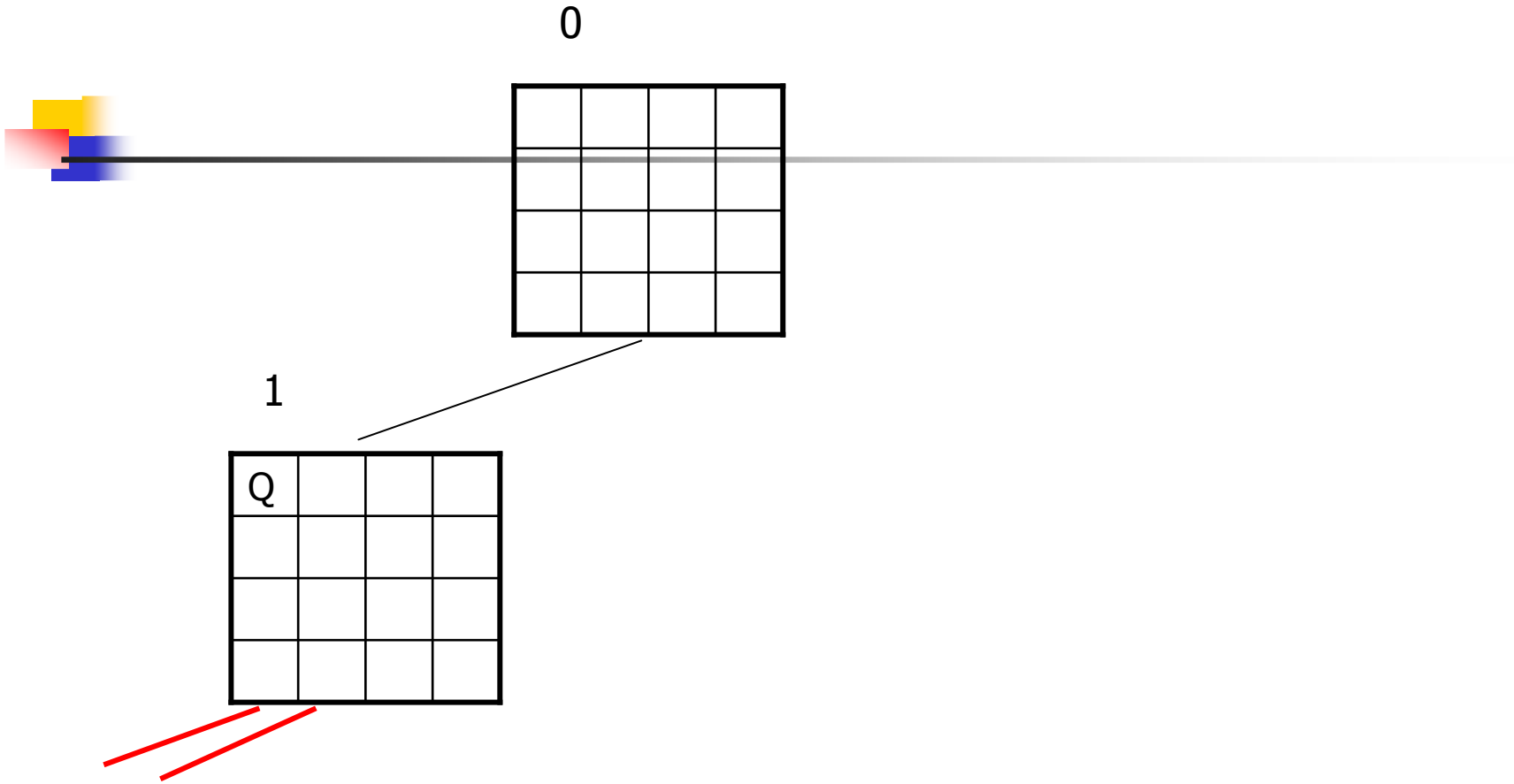
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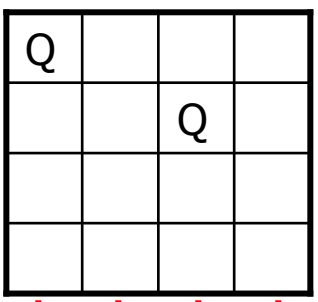
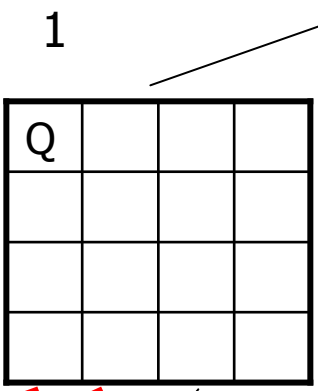
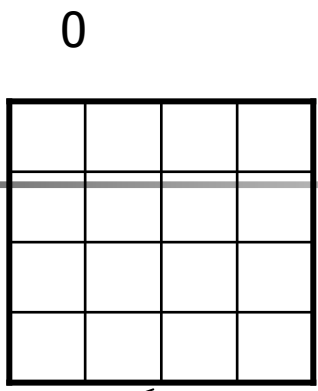
- Using Backtracking
  - *Large instances of difficult combinatorial problems can be solved*
  - *Worst case complexity of Backtracking can be exponential*
- Typically, a path is taken to check if a solution can be reached
  - *If not, the path is abandoned and another path taken*
  - *The process is repeated until the solution is arrived at*

# N-Queens problem

- Place  $n$ -queens on an  $n \times n$  chess board so that no two queens attack each other.
  - *A queen can attack another if the latter is on the same row, column or diagonal*

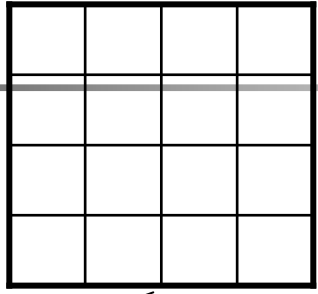




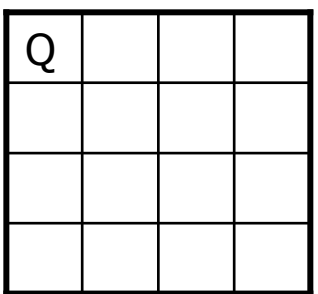




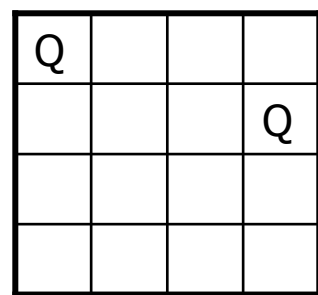
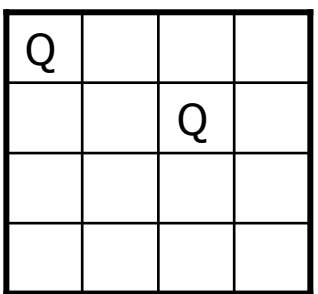
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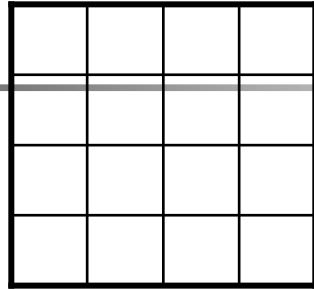


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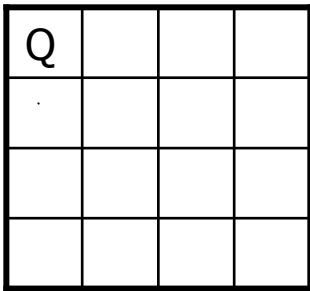
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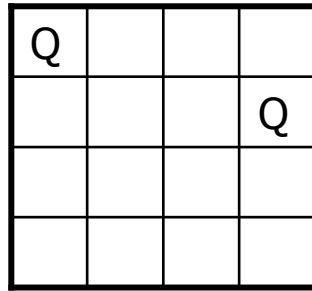
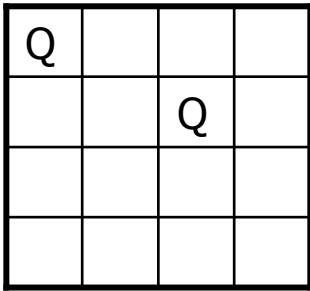
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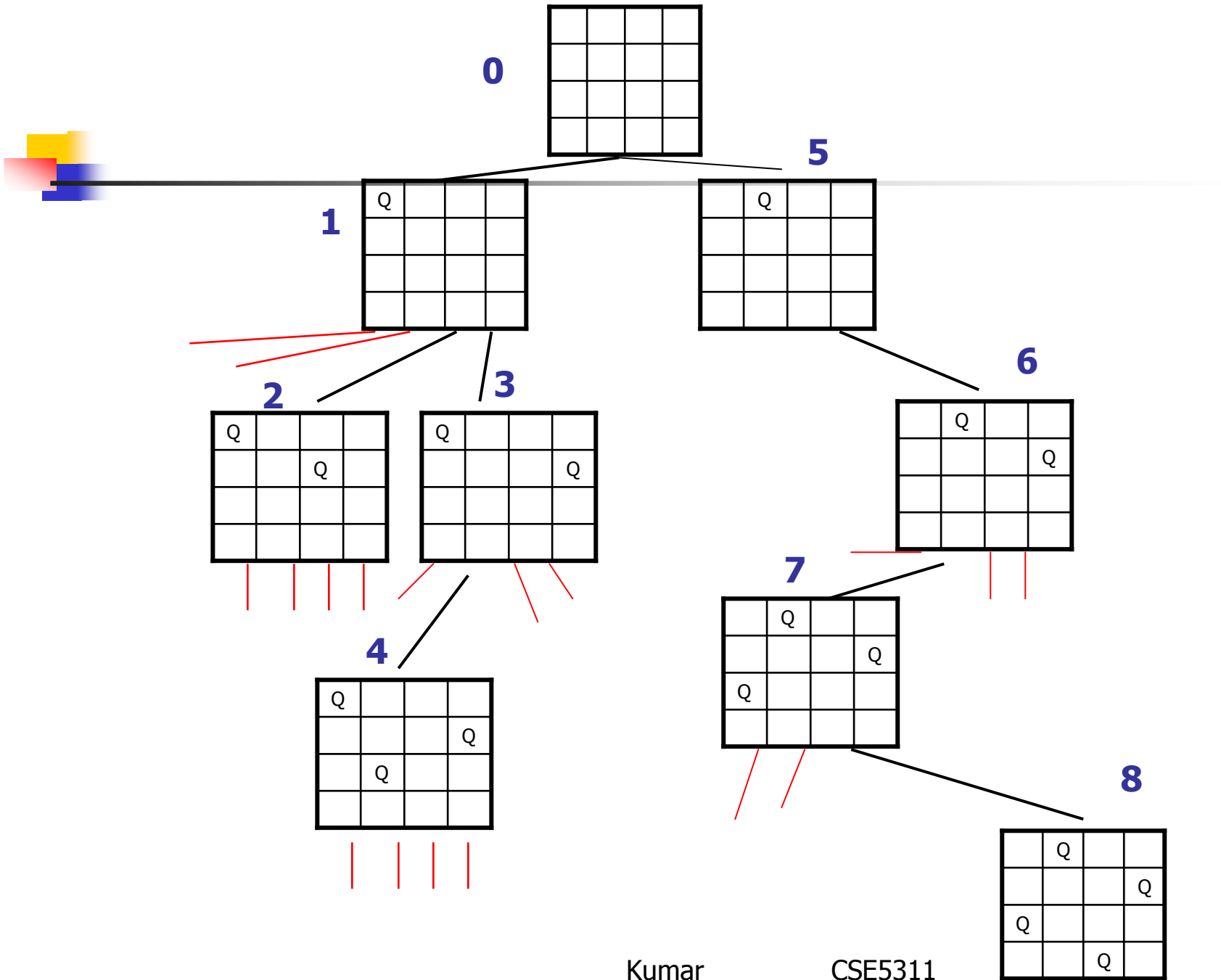


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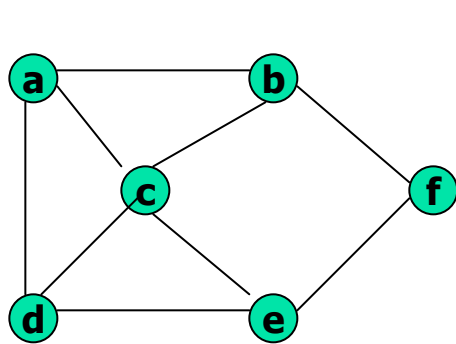
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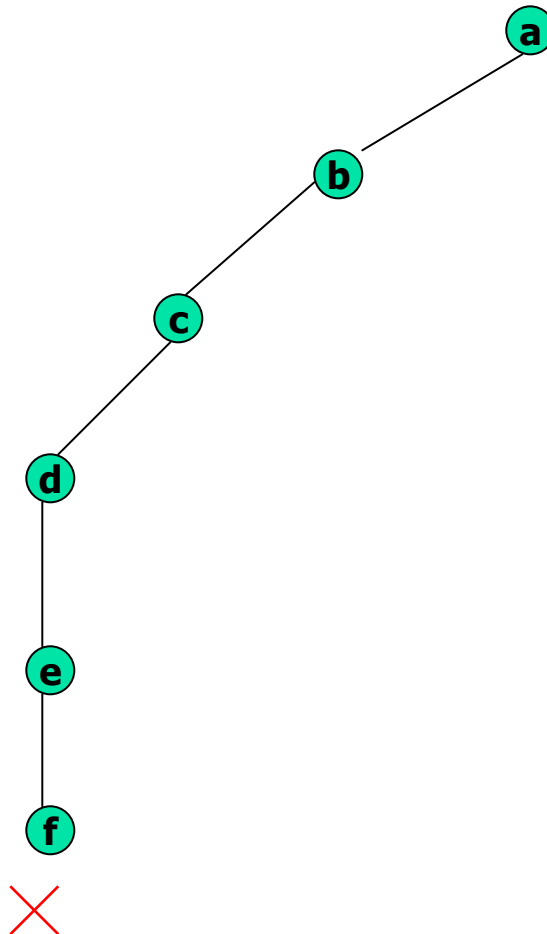




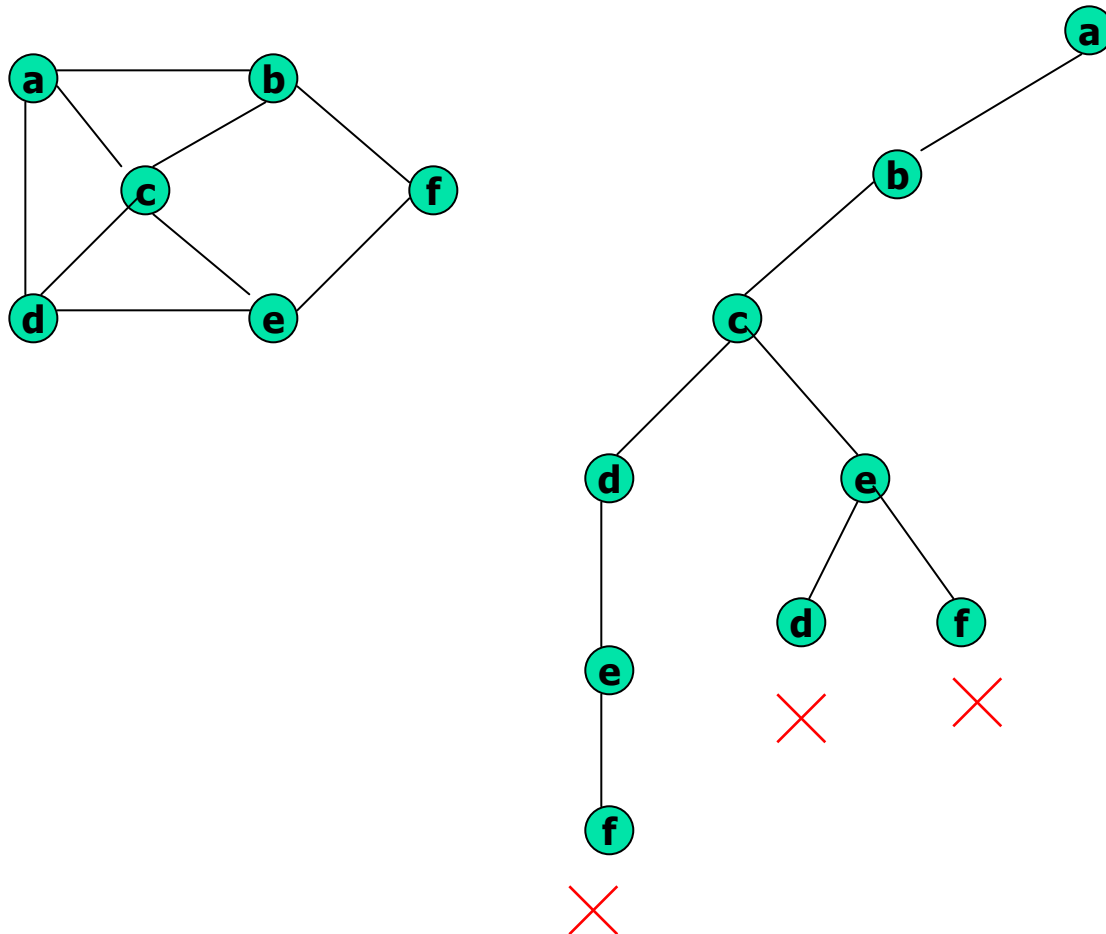
# Hamiltonian Circuit Problem



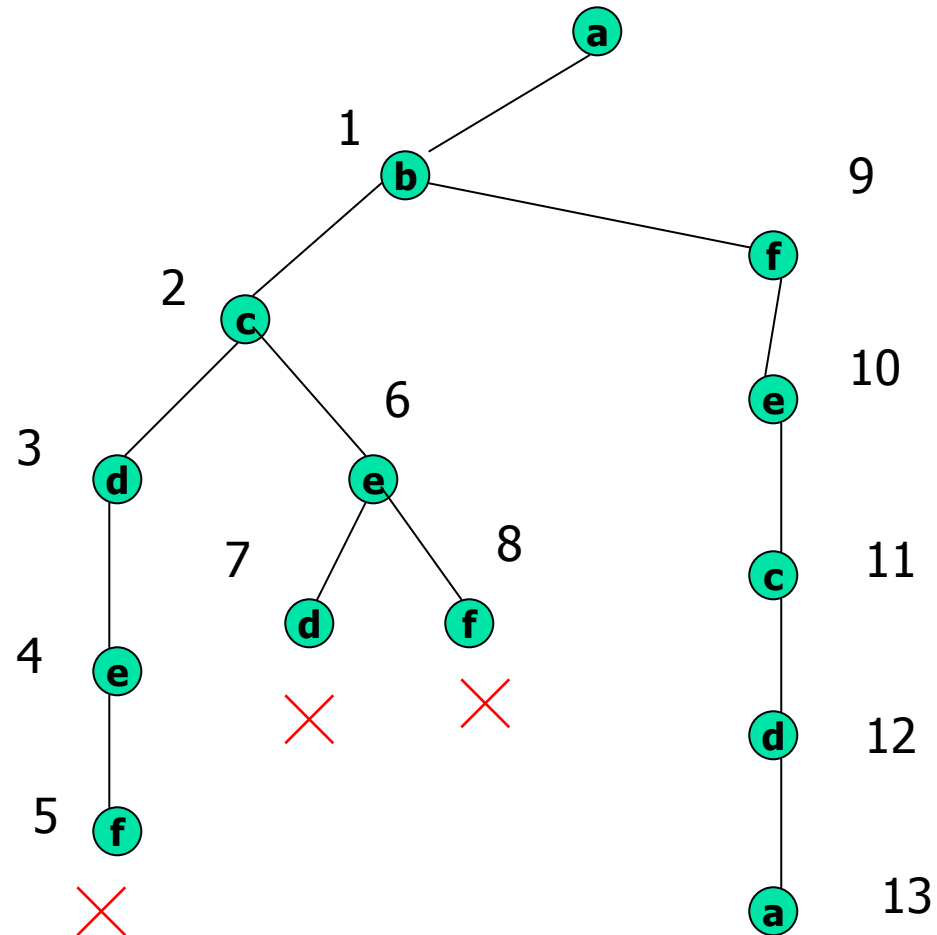
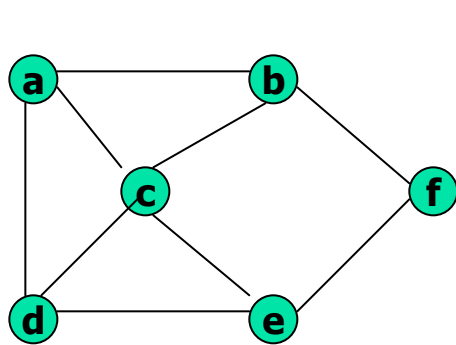
Start at a vertex and visit all the other vertices in the graph exactly once and return to the start vertex



# Hamiltonian Circuit Problem



# Hamiltonian Circuit Problem





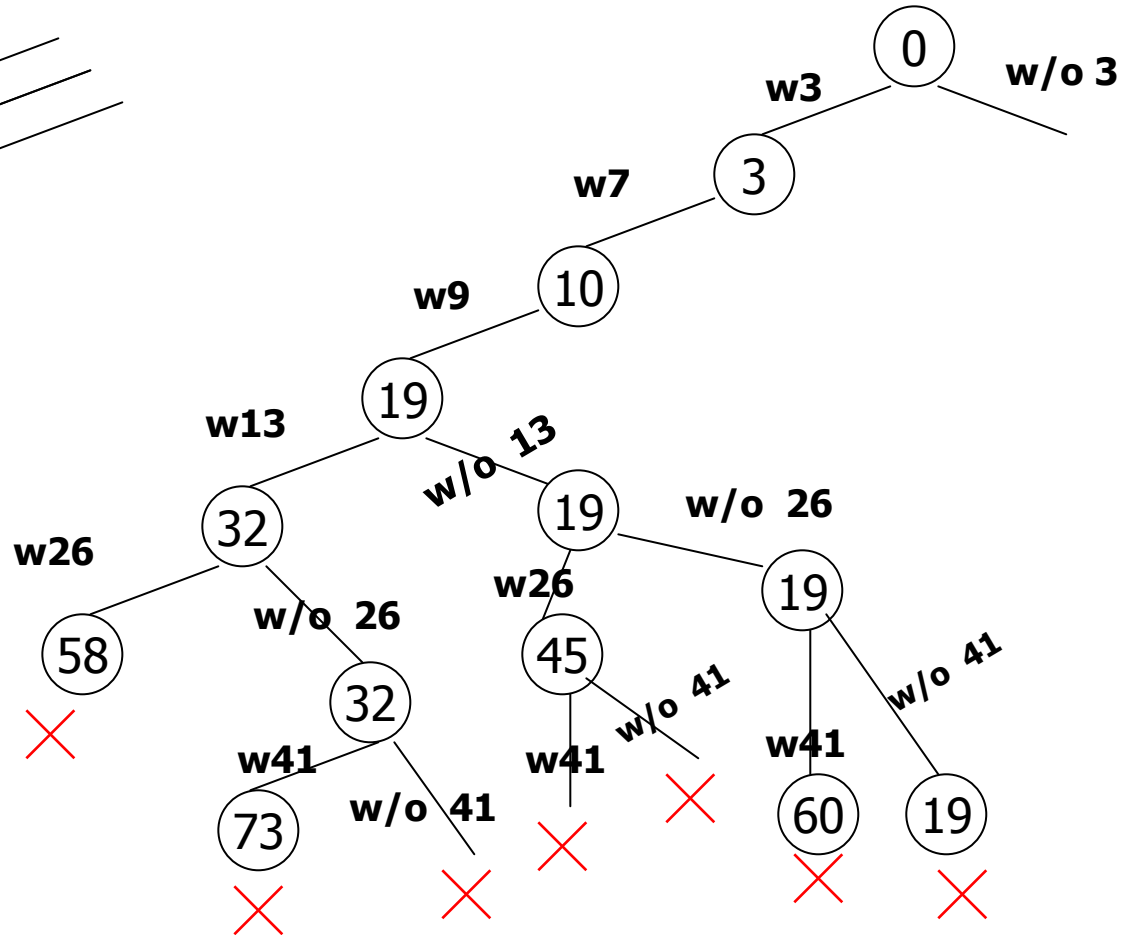
# Subset Sum Problem

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- Given a Set  $S = \{s_1, s_2, \dots, s_n\}$  and a positive integer 'd' find a subset of the given set S such that the sum of the positive integers in the subset is equal to 'd'.
- Let  $S = \{3, 7, 9, 13, 26, 41\}$ ;  $d = 51$ .
- Note – the list should be sorted.

# Subset problem

Let  $S = \{3, 7, 9, 13, 26, 41\}$ ;  
 $d = 51$







# Branch and Bound

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- With backtracking
  - *The search space is can be very large*
  - *It is an exhaustive search*
  - *Worst case complexity is exponential*
- Branch and bound technique
  - *Limits the search space*
  - *Through an estimate of the*
    - Upper bound or
    - Lower bound

# Scheduling problem

- The problem of assigning  $n$  people to  $n$  jobs such that the total cost is as small as possible

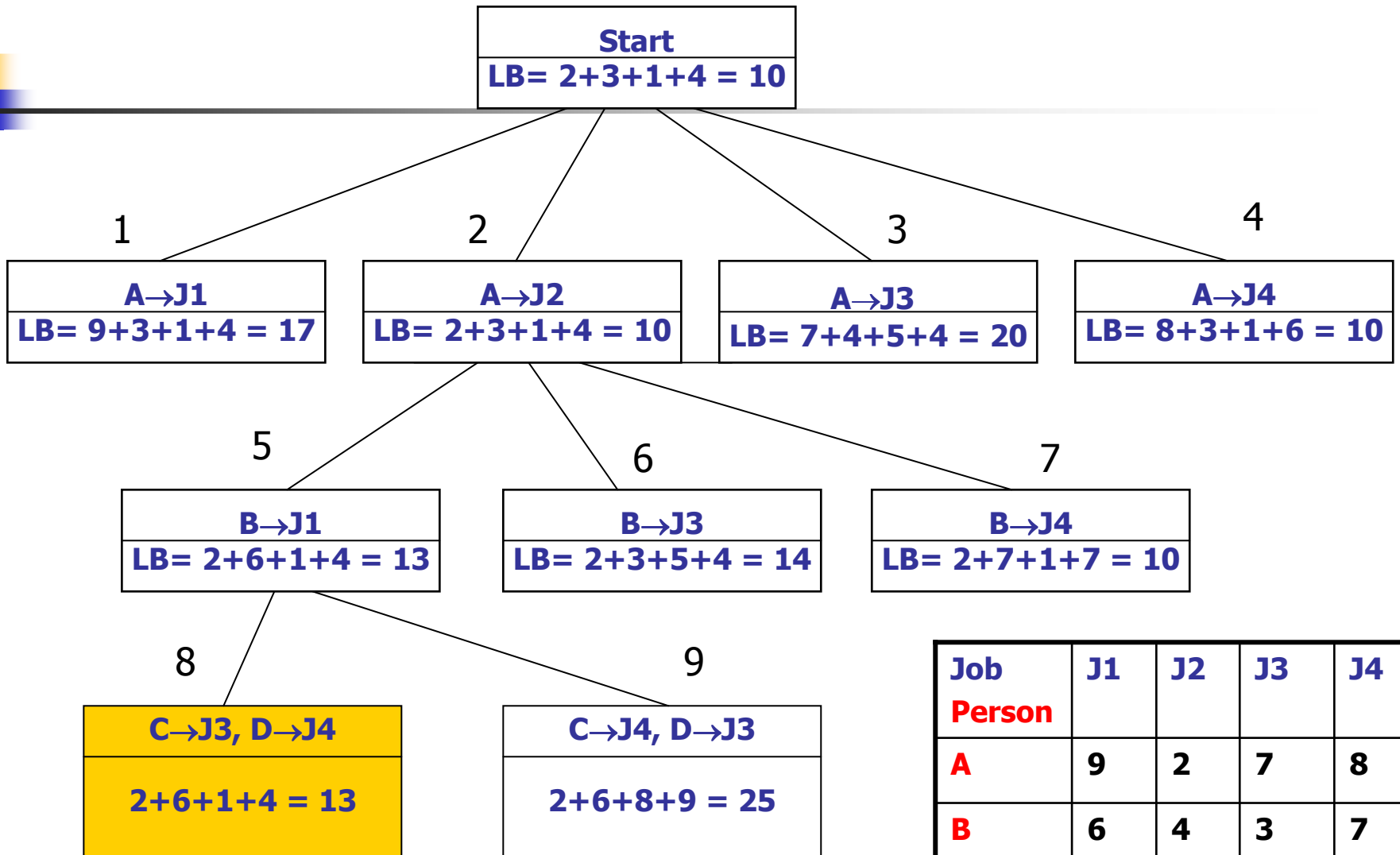
| <b>Job</b>    | <b>J1</b> | <b>J2</b> | <b>J3</b> | <b>J4</b> |
|---------------|-----------|-----------|-----------|-----------|
| <b>Person</b> |           |           |           |           |
| <b>A</b>      | 9         | 2         | 7         | 8         |
| <b>B</b>      | 6         | 4         | 3         | 7         |
| <b>C</b>      | 5         | 8         | 1         | 8         |
| <b>D</b>      | 7         | 6         | 9         | 4         |



# Branch and Bound

- Find a Lower Bound on the cost of the solution
- The lower bound is only an estimate
  - *This is only an estimate*
  - *The LB may not be a legitimate solution*
- In this case, consider the lowest cost form each row
  - $2 + 3 + 1 + 4 = 10$
  - *This is our LB*

| Job<br>Person | J1 | J2 | J3 | J4 |
|---------------|----|----|----|----|
| A             | 9  | 2  | 7  | 8  |
| B             | 6  | 4  | 3  | 7  |
| C             | 5  | 8  | 1  | 8  |
| D             | 7  | 6  | 9  | 4  |



| Job           | J1 | J2 | J3 | J4 |
|---------------|----|----|----|----|
| <b>Person</b> |    |    |    |    |
| <b>A</b>      | 9  | 2  | 7  | 8  |
| <b>B</b>      | 6  | 4  | 3  | 7  |
| <b>C</b>      | 5  | 8  | 1  | 8  |
| <b>D</b>      | 7  | 6  | 9  | 4  |

# Knapsack Problem

- We wish to maximize the profit in the knapsack
- Maximization
- Use Upper bound
- $UB = v + (W-v)(v_{i+1}/w_{i+1})$
- When we start  $v = 0$

**W = 10**

| Item     | Weight   | Value       | Value/<br>weight |  |
|----------|----------|-------------|------------------|--|
| <b>1</b> | <b>4</b> | <b>\$40</b> | <b>10</b>        |  |
| <b>2</b> | <b>7</b> | <b>\$42</b> | <b>6</b>         |  |
| <b>3</b> | <b>5</b> | <b>\$25</b> | <b>5</b>         |  |
| <b>4</b> | <b>3</b> | <b>\$12</b> | <b>4</b>         |  |

# Traveling Salesperson Problem

- $LB = \sum (\text{distance to two nearest cities})/2$
- $\sum$  over all cities

