

Shortest Paths

CSE 2320 – Algorithms and Data Structures
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Terminology

- A **network** is a **directed graph**. We will use both terms interchangeably.
- The **weight of a path** is the sum of weights of the edges that make up the path.
- The **shortest path** between two vertices s and t in a directed graph is a directed path from s to t with the property that no other such path has a lower weight.

Shortest Paths

- Finding shortest paths is not a single problem, but rather a family of problems.
- We will consider two of these problems:
 - Single-source: find the shortest path from the source vertex v to all other vertices in the graph.
 - It turns out that these shortest paths form a tree, with v as the root.
 - All-pairs: find the shortest paths for all pairs of vertices in the graph.

Assumptions

- Allow directed graphs.
 - In all our shortest path algorithms, we will allow graphs to be directed.
 - Obviously, any algorithm that works on directed graphs will also work on undirected graphs. Why?
- Negative edge weights are not allowed. Why?

Assumptions

- Allow directed graphs.
 - In all our shortest path algorithms, we will allow graphs to be directed.
 - Obviously, any algorithm that works on directed graphs will also work on undirected graphs. Why?
 - Undirected graphs are a special case of directed graphs.
- Negative edge weights are not allowed. Why?
 - With negative weights, "shortest paths" may not be defined.
 - If a cyclic path has negative weight, then repeating that path infinitely will lead to "shorter" and "shorter" paths.
 - If all weights are nonnegative, a shortest path never needs to include a cycle.

Shortest-Paths Spanning Tree

- Given a network G and a designated vertex s , a **shortest-paths spanning tree** (SPST) for s is a tree that contains s and all vertices reachable from s , such that:
 - Vertex s is the root of this tree.
 - Each tree path is a shortest path in G .

Computing SPSTs

- To compute an SPST, given a graph G and a vertex s , we will design an algorithm that maintains and updates the following two arrays:
 - Array wt : $wt[v]$ is the weight of the shortest path we have found so far from s to v .
 - At the beginning, $wt[v] = \text{infinity}$, except for s , where $wt[s] = 0$.
 - Array st : $st[v]$ is the parent vertex of v on the shortest path found so far from s to v .
 - At the beginning, $st[v] = -1$, except for s , where $st[s] = s$.
 - Array in : $in[v]$ is 1 if v has been already added to the SPST, 0 otherwise.
 - At the beginning, $in[v] = 0$, except for s , where $in[s] = 1$.

Dijkstra's Algorithm

- Computes an SPST for a graph G and a source s .
- Very similar to Prim's algorithm, but:
 - First vertex to add is the source.
 - Works with directed graphs, whereas Prim's only works with undirected graphs.
 - Requires edge weights to be non-negative.
 - **The wt array behaves differently (see next slides).**
- Time: $O(V^2)$, similar analysis to that of Prim's algorithm.
- Time $O(E \lg V)$ using a priority-queue implementation.

Dijkstra's Algorithm

Input: number of vertices V , $V \times V$ array weight, source vertex s .

1. For all v :
 2. $wt[v] = \text{infinity}$.
 3. $st[v] = -1$.
 4. $in[v] = 0$.
5. $wt[s] = 0, st[s] = s$.
6. Repeat until all vertices have been added to the tree:
 7. Find the v with the smallest $wt[v]$, among all v such that $in[v] = 0$.
 8. Add to the SPST vertex v and edge from $st[v]$ to v .
 9. $in[v] = 1$.
 10. For each neighbor w of v , such that $in[w] = 0$:
 - 11. If $wt[w] > wt[v] + \text{weight}[v, w]$:**
 - 12. $wt[w] = wt[v] + \text{weight}[v, w]$,**
 - 13. $st[w] = v$.**

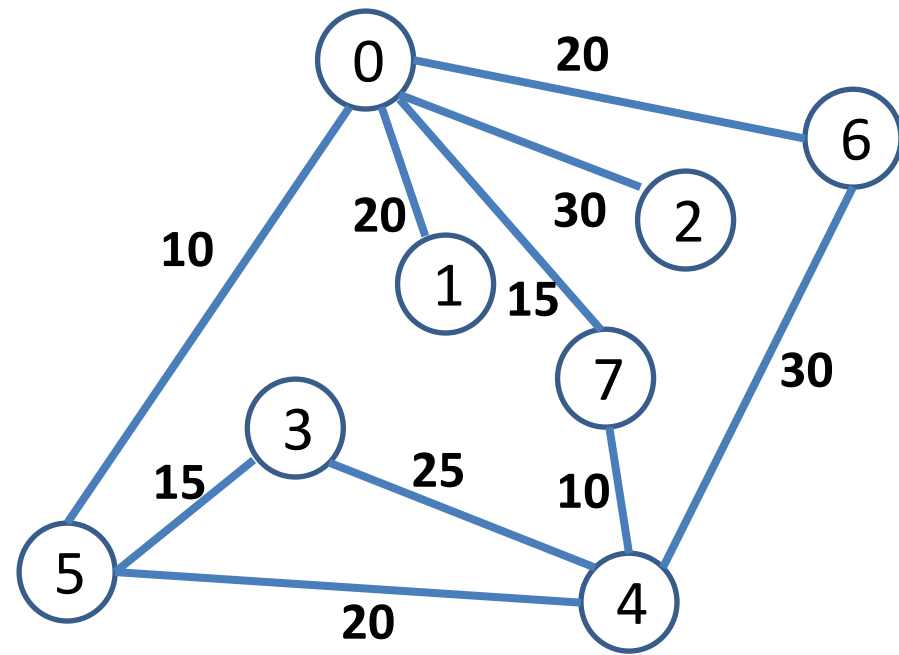
Edge Relaxation

```
if (wt[w] > wt[v] + e.wt)
{
    wt[w] = wt[v] + e.wt;
    st[w] = v;
}
```

- $wt[w]$: current estimate of shortest distance from source to w .
- $st[w]$: parent vertex of w on shortest found path from source to w .

Dijkstra Example

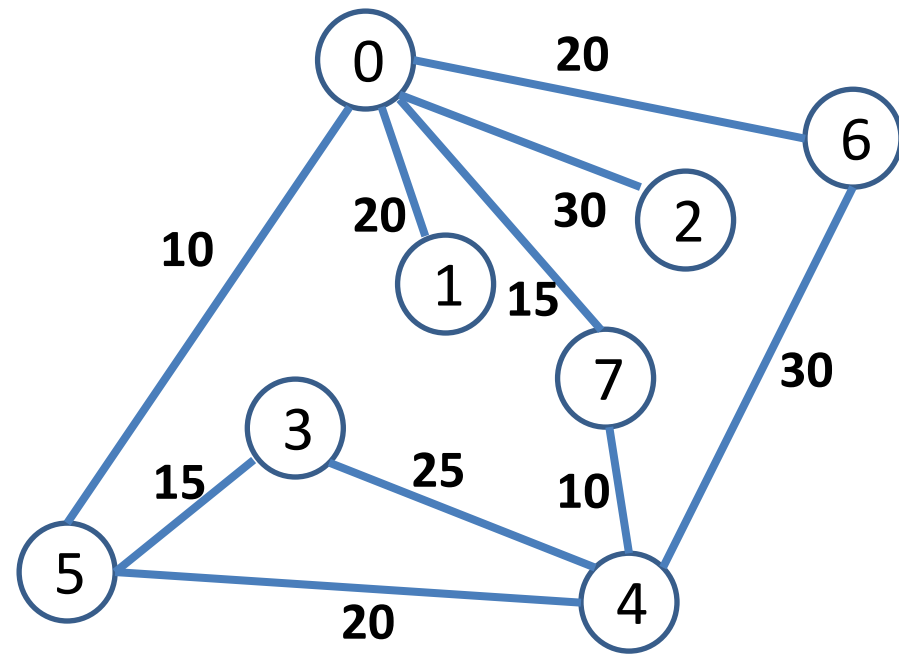
- Suppose we want to compute the SPST for vertex 7.
- First, we initialize arrays wt, st, in (steps 2, 3, 4).



vertex	0	1	2	3	4	5	6	7
wt	inf	inf	inf	inf	inf	inf	inf	inf
st	-1	-1	-1	-1	-1	-1	-1	-1
in	0	0	0	0	0	0	0	0

Dijkstra Example

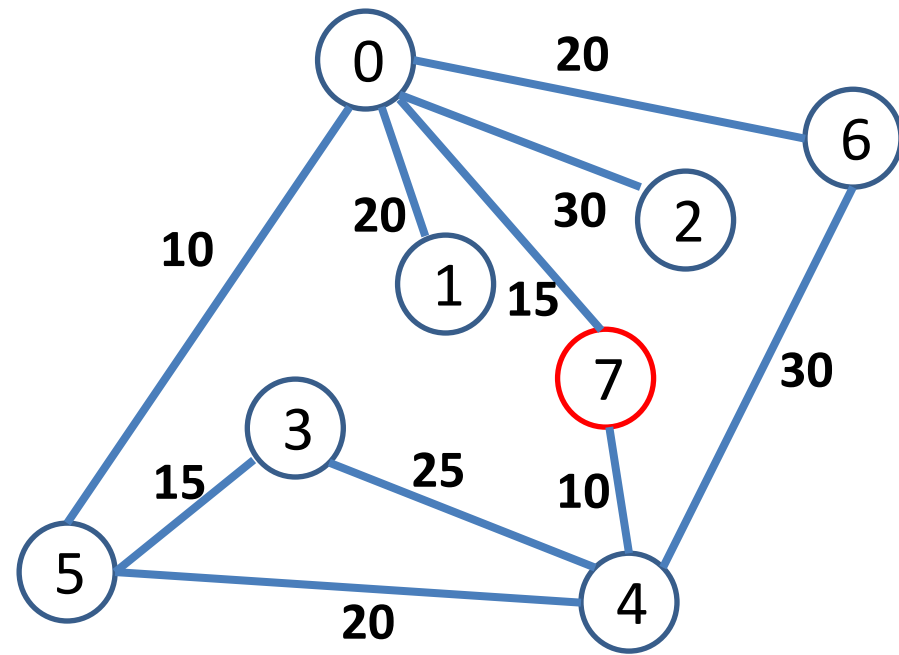
- Suppose we want to compute the SPST for vertex 7.
- Step 5.



vertex	0	1	2	3	4	5	6	7
wt	inf	inf	inf	inf	inf	inf	inf	0
st	-1	-1	-1	-1	-1	-1	-1	7
in	0	0	0	0	0	0	0	0

Dijkstra Example

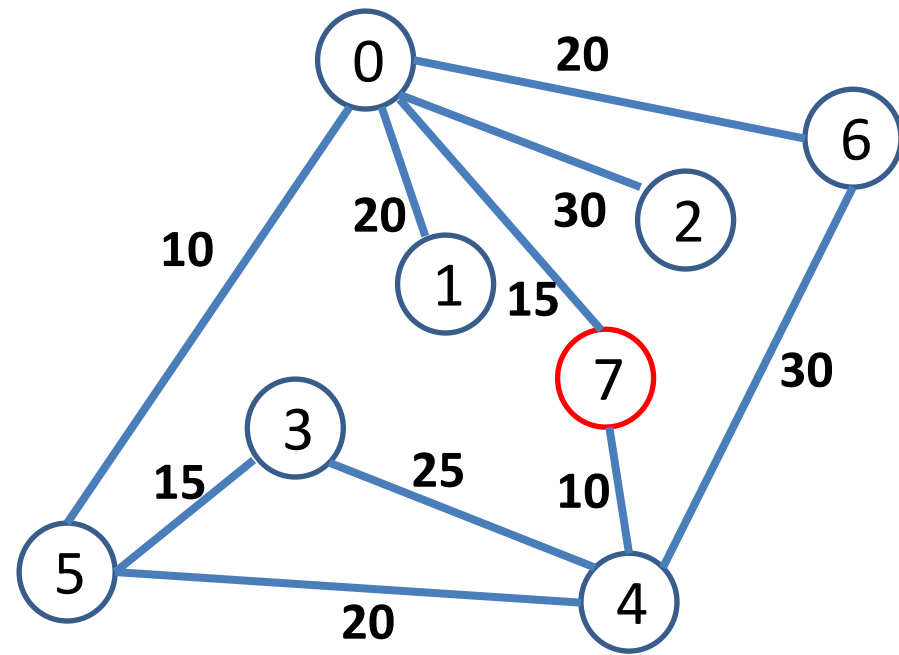
- Suppose we want to compute the SPST for vertex 7.
- Steps 7, 8, 9: $v = 7$



vertex	0	1	2	3	4	5	6	7
wt	inf	inf	inf	inf	inf	inf	inf	0
st	-1	-1	-1	-1	-1	-1	-1	7
in	0	0	0	0	0	0	0	1

Dijkstra Example

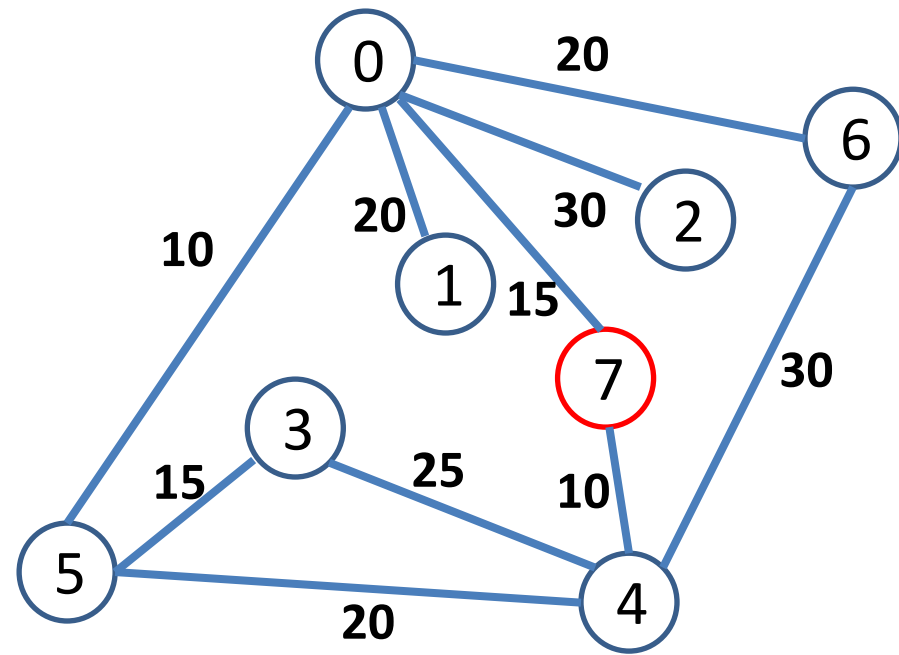
- Suppose we want to compute the SPST for vertex 7.
- Steps 7, 8, 9: $v = 7$
- Step 10: For $w = \{0, 4\}$
 - Step 11: Compare inf with 15
 - Steps 12, 13: $wt[0] = wt[7] + 15$, $st[0] = 7$.



vertex	0	1	2	3	4	5	6	7
wt	15	inf	inf	inf	inf	inf	inf	0
st	7	-1	-1	-1	-1	-1	-1	7
in	0	0	0	0	0	0	0	1

Dijkstra Example

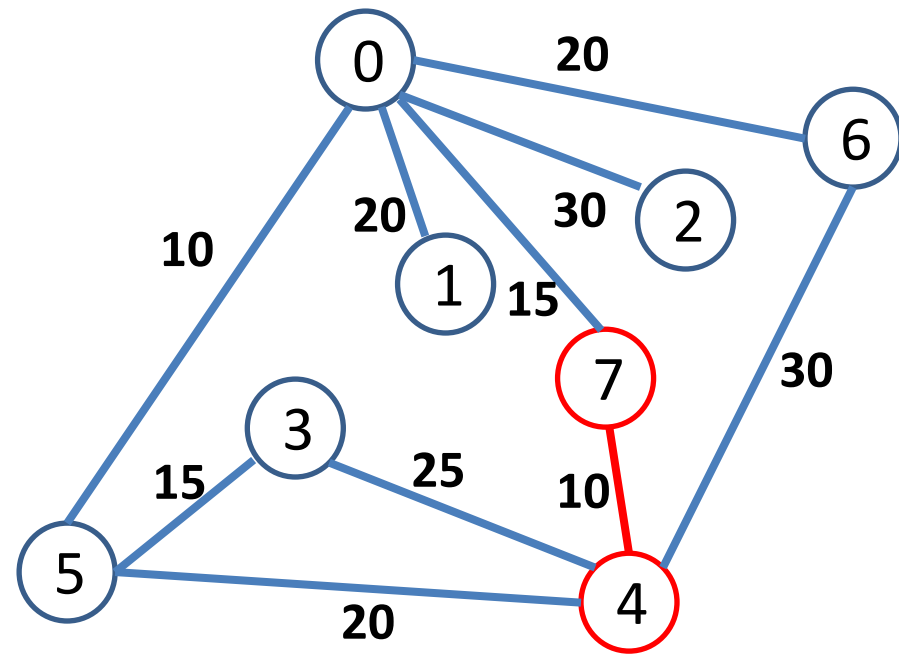
- Suppose we want to compute the SPST for vertex 7.
- Steps 7, 8, 9: $v = 7$
- Step 10: For $w = \{0, 4\}$
 - Step 11: Compare inf with 10
 - Steps 12, 13: $wt[4] = wt[7] + 10$, $st[4] = 7$.



vertex	0	1	2	3	4	5	6	7
wt	15	inf	inf	inf	10	inf	inf	0
st	7	-1	-1	-1	7	-1	-1	7
in	0	0	0	0	0	0	0	1

Dijkstra Example

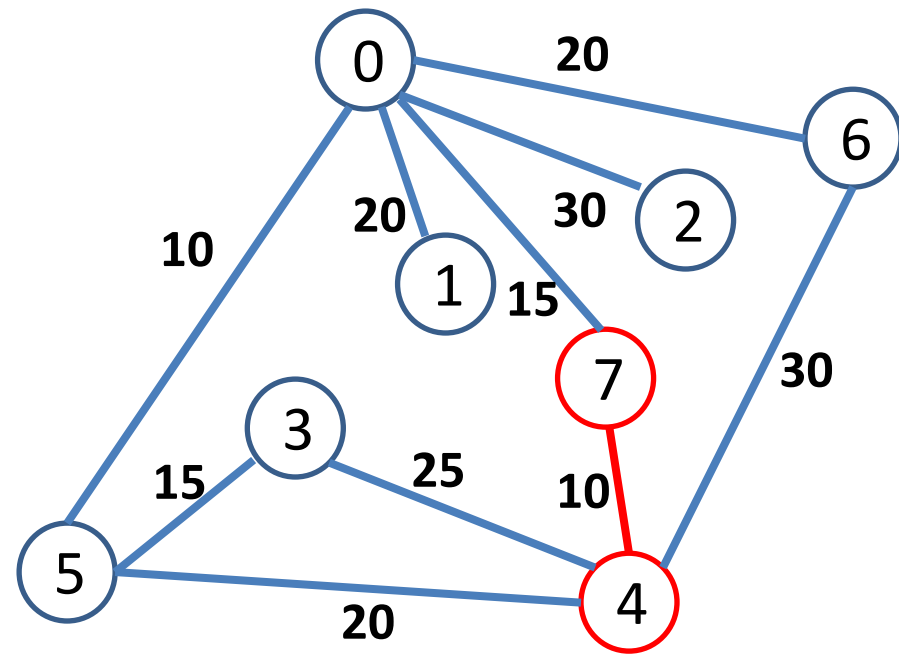
- Suppose we want to compute the SPST for vertex 7.
- Steps 7, 8, 9: $v = 4$



vertex	0	1	2	3	4	5	6	7
wt	15	inf	inf	inf	10	inf	inf	0
st	7	-1	-1	-1	7	-1	-1	7
in	0	0	0	0	1	0	0	1

Dijkstra Example

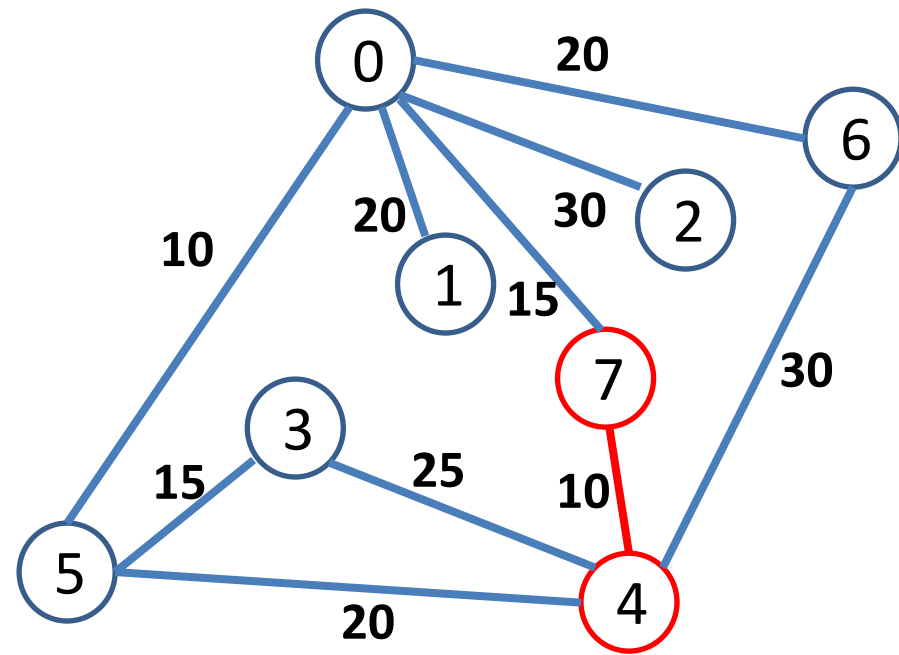
- Suppose we want to compute the SPST for vertex 7.
- Steps 7, 8, 9: $v = 4$
- Step 10: For $w = \{3, 5, 6\}$
 - Step 11: Compare inf with $10 + 25 = 35$
 - Steps 12, 13: $wt[3] = wt[4] + 25$, $st[3] = 4$.



vertex	0	1	2	3	4	5	6	7
wt	15	inf	inf	35	10	inf	inf	0
st	7	-1	-1	4	7	-1	-1	7
in	0	0	0	0	1	0	0	1

Dijkstra Example

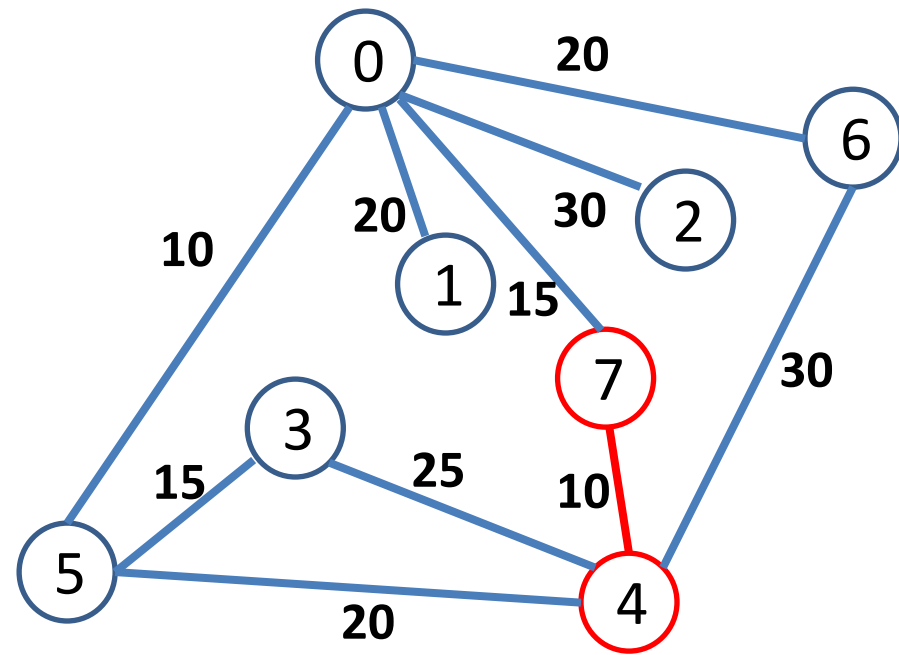
- Suppose we want to compute the SPST for vertex 7.
- Steps 7, 8, 9: $v = 4$
- Step 10: For $w = \{3, 5, 6\}$
 - Step 11: Compare inf with $10 + 20 = 30$
 - Steps 12, 13: $\text{wt}[5] = \text{wt}[4] + 20$, $\text{st}[5] = 4$.



vertex	0	1	2	3	4	5	6	7
wt	15	inf	inf	35	10	30	inf	0
st	7	-1	-1	4	7	4	-1	7
in	0	0	0	0	1	0	0	1

Dijkstra Example

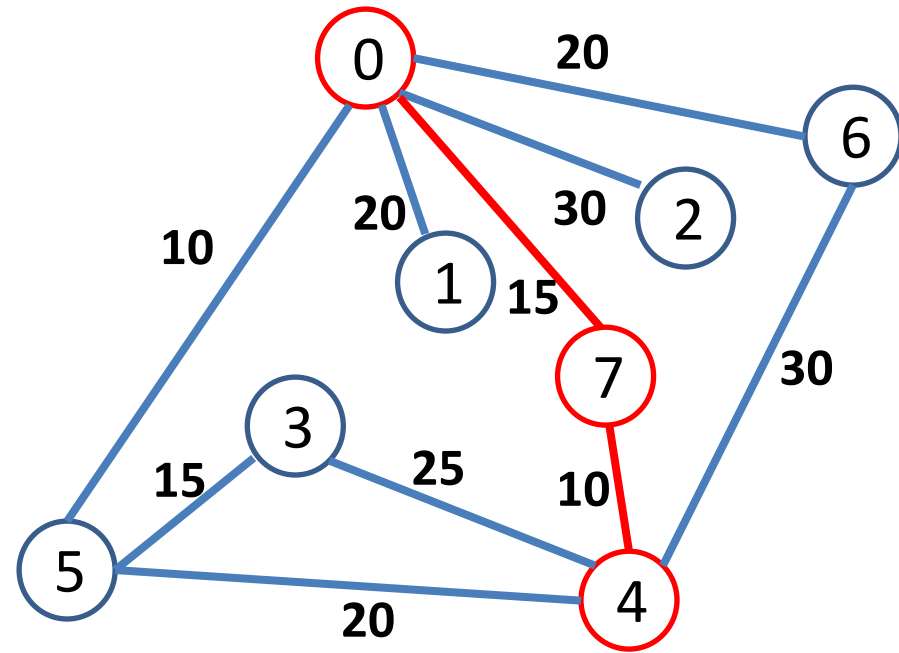
- Suppose we want to compute the SPST for vertex 7.
- Steps 7, 8, 9: $v = 4$
- Step 10: For $w = \{3, 5, 6\}$
 - Step 11: Compare inf with $10+30=40$
 - Steps 12, 13: $wt[6] = wt[4] + 30$, $st[6] = 4$.



vertex	0	1	2	3	4	5	6	7
wt	15	inf	inf	35	10	30	40	0
st	7	-1	-1	4	7	4	4	7
in	0	0	0	0	1	0	0	1

Dijkstra Example

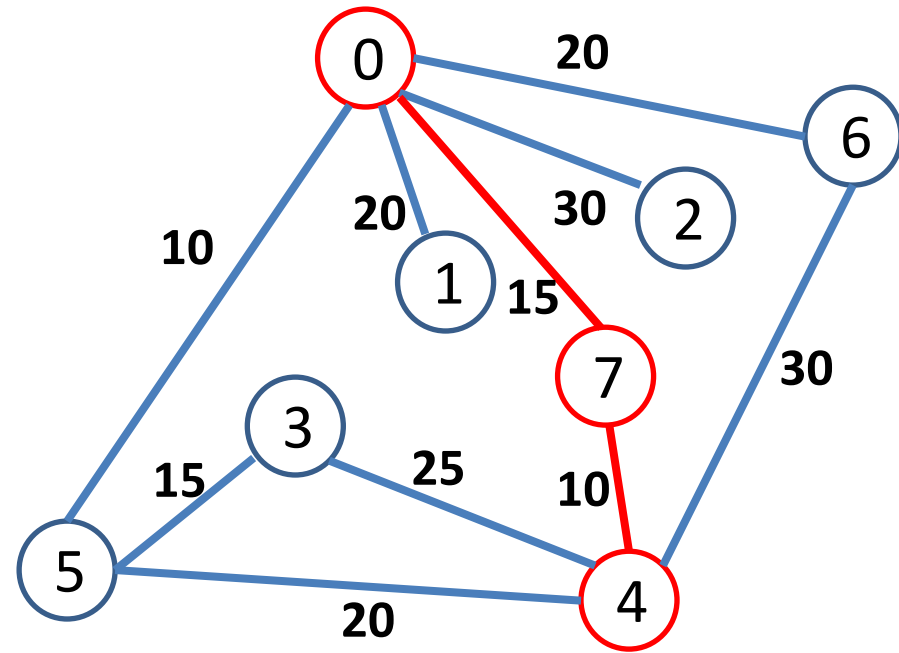
- Suppose we want to compute the SPST for vertex 7.
- Steps 7, 8, 9: $v = 0$



vertex	0	1	2	3	4	5	6	7
wt	15	inf	inf	35	10	30	40	0
st	7	-1	-1	4	7	4	4	7
in	1	0	0	0	1	0	0	1

Dijkstra Example

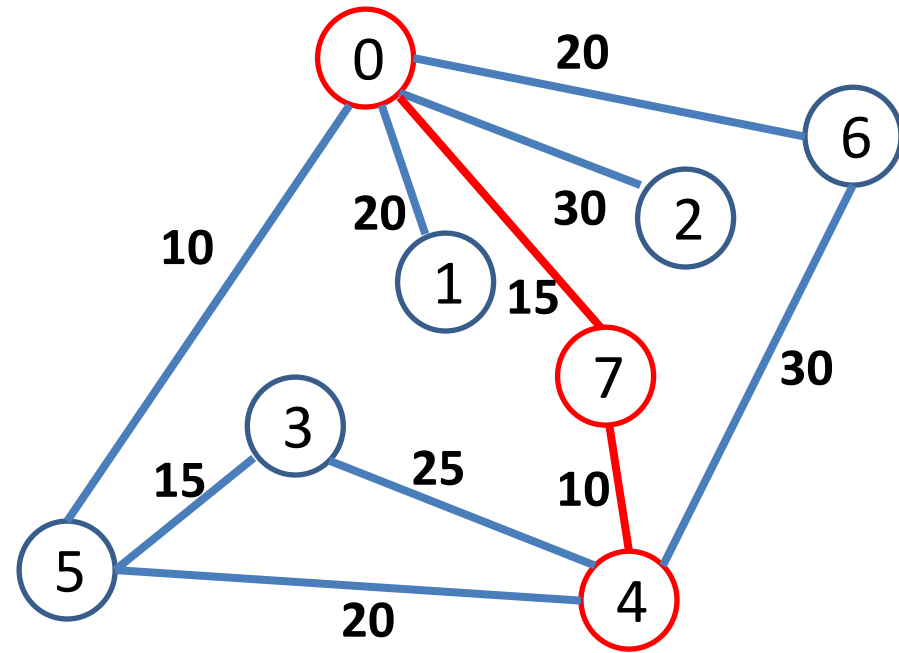
- Suppose we want to compute the SPST for vertex 7.
- Steps 7, 8, 9: $v = 0$
- Step 10: For $w = \{1, 2, 5, 6\}$
 - Step 11: Compare inf with $15+20=35$
 - Steps 12, 13: $wt[1] = wt[0] + 20$, $st[1] = 0$.



vertex	0	1	2	3	4	5	6	7
wt	15	35	inf	35	10	30	40	0
st	7	0	-1	4	7	4	4	7
in	1	0	0	0	1	0	0	1

Dijkstra Example

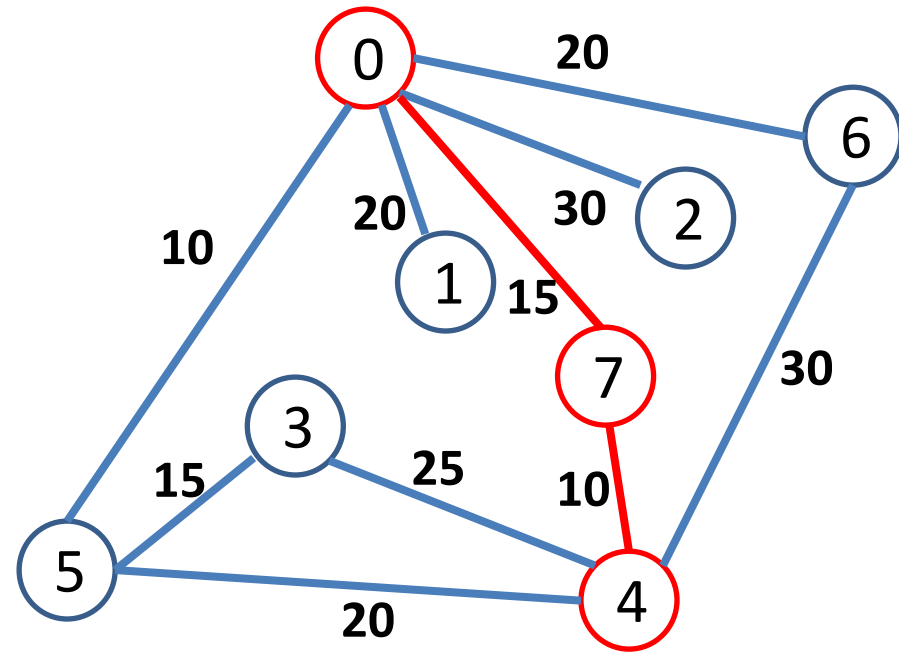
- Suppose we want to compute the SPST for vertex 7.
- Steps 7, 8, 9: $v = 0$
- Step 10: For $w = \{1, 2, 5, 6\}$
 - Step 11: Compare inf with $15+30=45$
 - Steps 12, 13: $wt[2] = wt[0] + 30$, $st[2] = 0$.



vertex	0	1	2	3	4	5	6	7
wt	15	35	45	35	10	30	40	0
st	7	0	0	4	7	4	4	7
in	1	0	0	0	1	0	0	1

Dijkstra Example

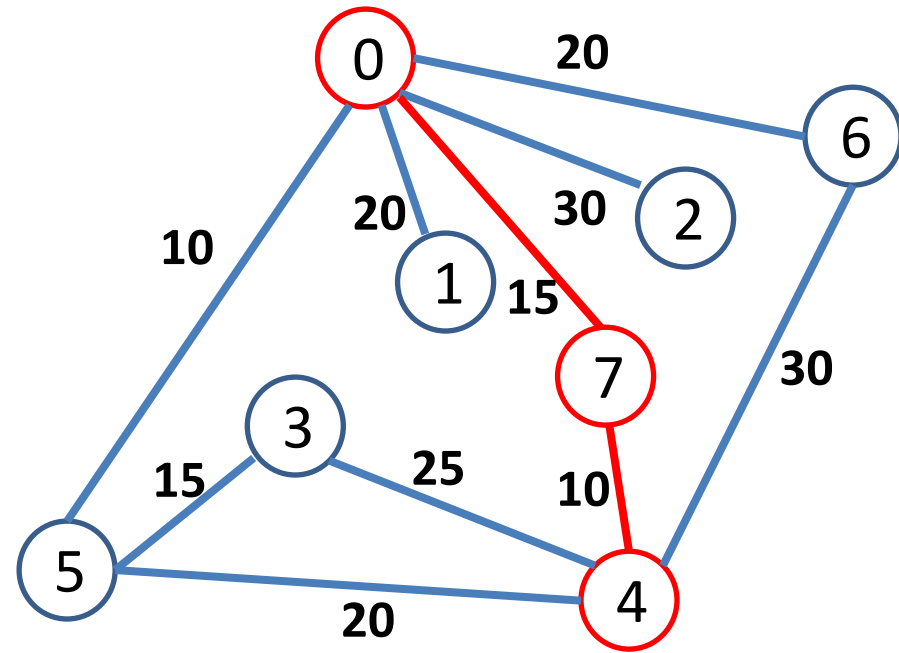
- Suppose we want to compute the SPST for vertex 7.
- Steps 7, 8, 9: $v = 0$
- Step 10: For $w = \{1, 2, 5, 6\}$
 - Step 11: Compare 30 with $15 + 10 = 25$
 - Steps 12, 13: $wt[5] = wt[0] + 10$, $st[5] = 0$.



vertex	0	1	2	3	4	5	6	7
wt	15	35	45	35	10	25	40	0
st	7	0	0	4	7	0	4	7
in	1	0	0	0	1	0	0	1

Dijkstra Example

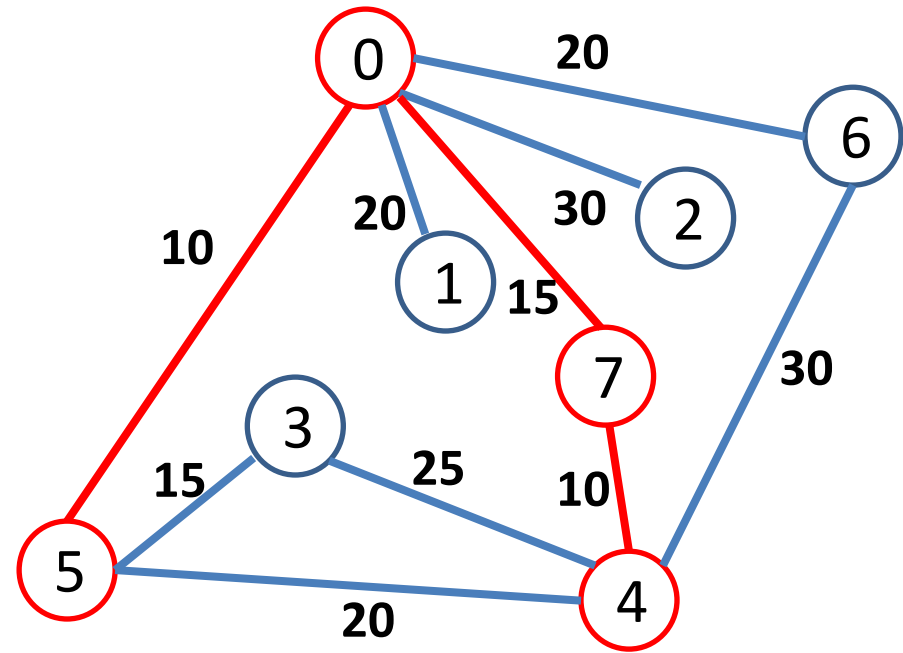
- Suppose we want to compute the SPST for vertex 7.
- Steps 7, 8, 9: $v = 0$
- Step 10: For $w = \{1, 2, 5, 6\}$
 - Step 11: Compare 40 with $15+20=35$
 - Steps 12, 13: $wt[6] = wt[0] + 20$, $st[6] = 0$.



vertex	0	1	2	3	4	5	6	7
wt	15	35	45	35	10	25	35	0
st	7	0	0	4	7	0	0	7
in	1	0	0	0	1	0	0	1

Dijkstra Example

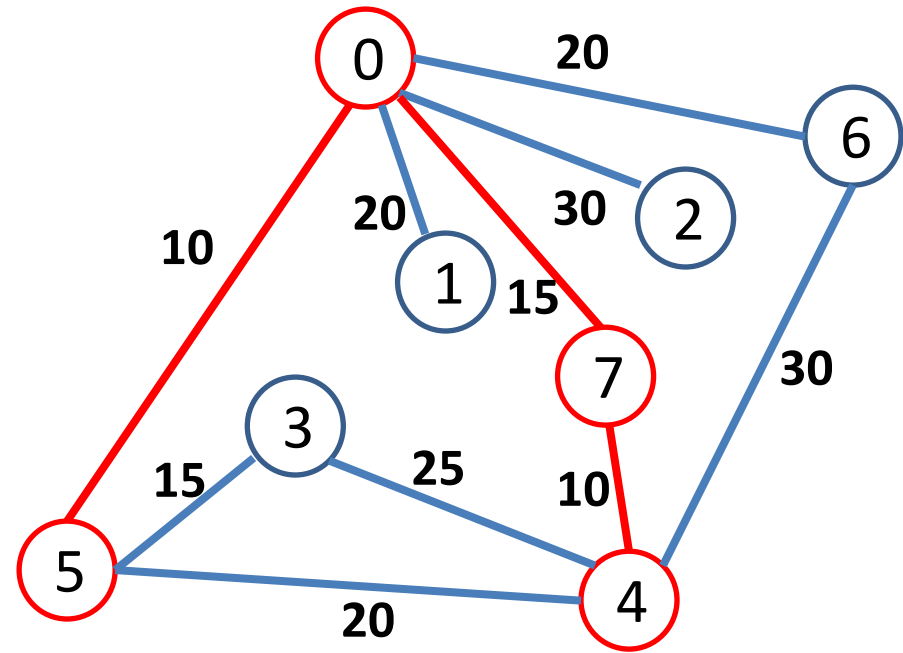
- Suppose we want to compute the SPST for vertex 7.
- Steps 7, 8, 9: $v = 5$



vertex	0	1	2	3	4	5	6	7
wt	15	35	45	35	10	25	35	0
st	7	0	0	4	7	0	0	7
in	1	0	0	0	1	1	0	1

Dijkstra Example

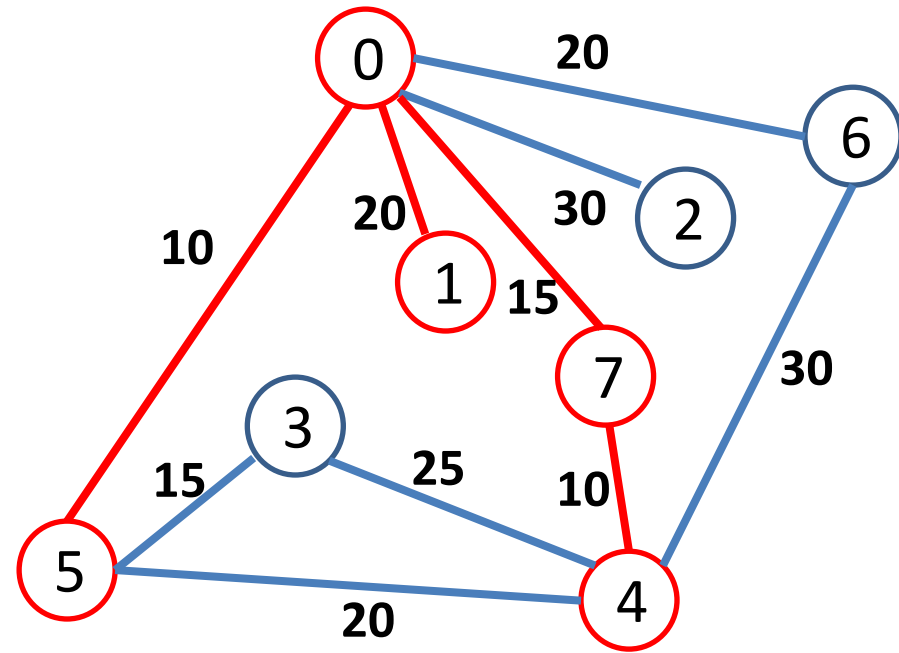
- Suppose we want to compute the SPST for vertex 7.
- Steps 7, 8, 9: $v = 5$
- Step 10: For $w = \{3\}$
 - Step 11: Compare 35 with $25 + 15 = 40$
NO UPDATE



vertex	0	1	2	3	4	5	6	7
wt	15	35	45	35	10	25	35	0
st	7	0	0	4	7	0	0	7
in	1	0	0	0	1	1	0	1

Dijkstra Example

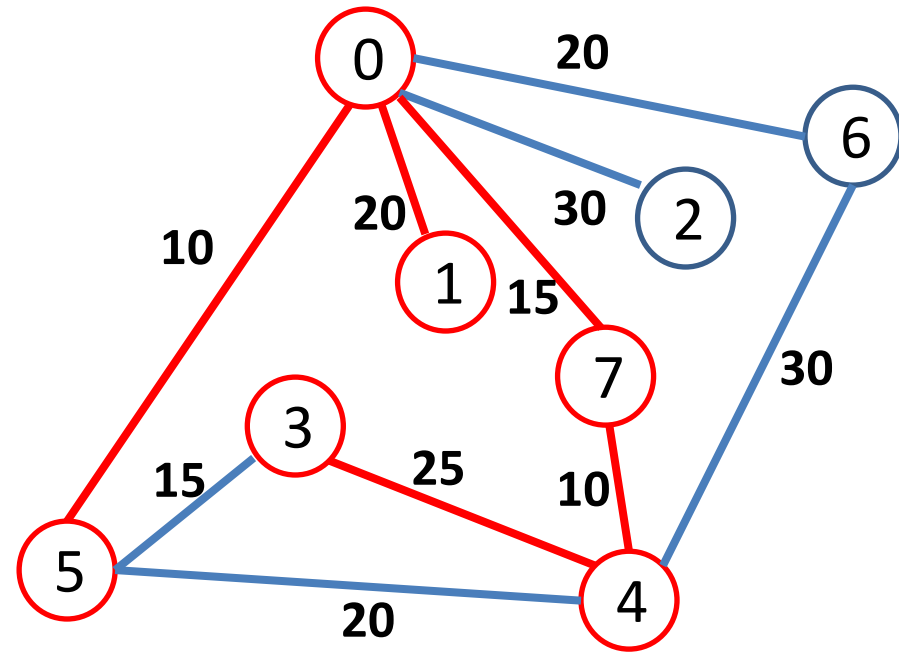
- Suppose we want to compute the SPST for vertex 7.
- Steps 7, 8, 9: $v = 1$
- Step 10: For $w =$ empty list



vertex	0	1	2	3	4	5	6	7
wt	15	35	45	35	10	25	35	0
st	7	0	0	4	7	0	0	7
in	1	1	0	0	1	1	0	1

Dijkstra Example

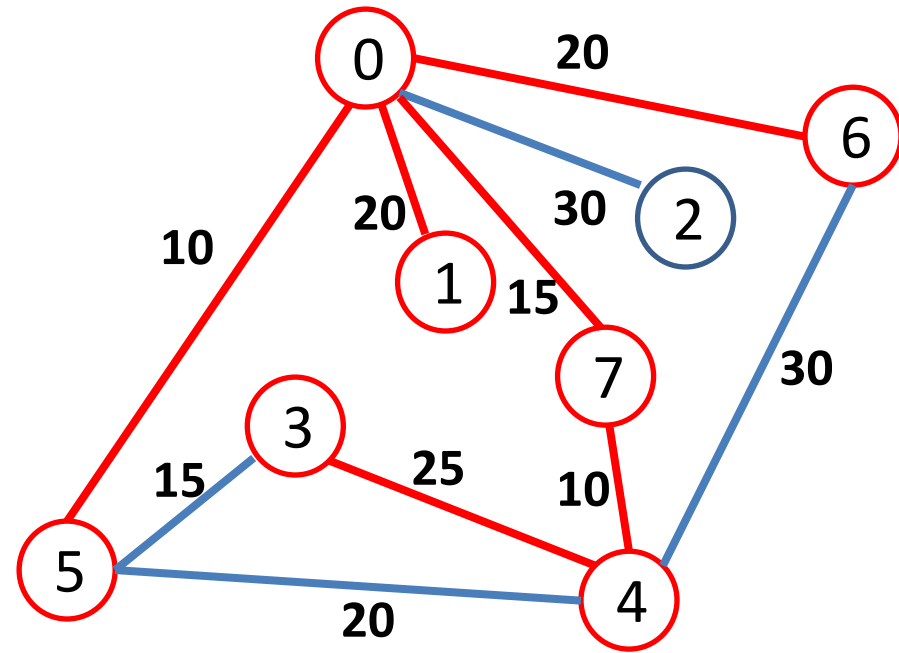
- Suppose we want to compute the SPST for vertex 7.
- Steps 7, 8, 9: $v = 3$
- Step 10: For $w =$ empty list



vertex	0	1	2	3	4	5	6	7
wt	15	35	45	35	10	25	35	0
st	7	0	0	4	7	0	0	7
in	1	1	0	1	1	1	0	1

Dijkstra Example

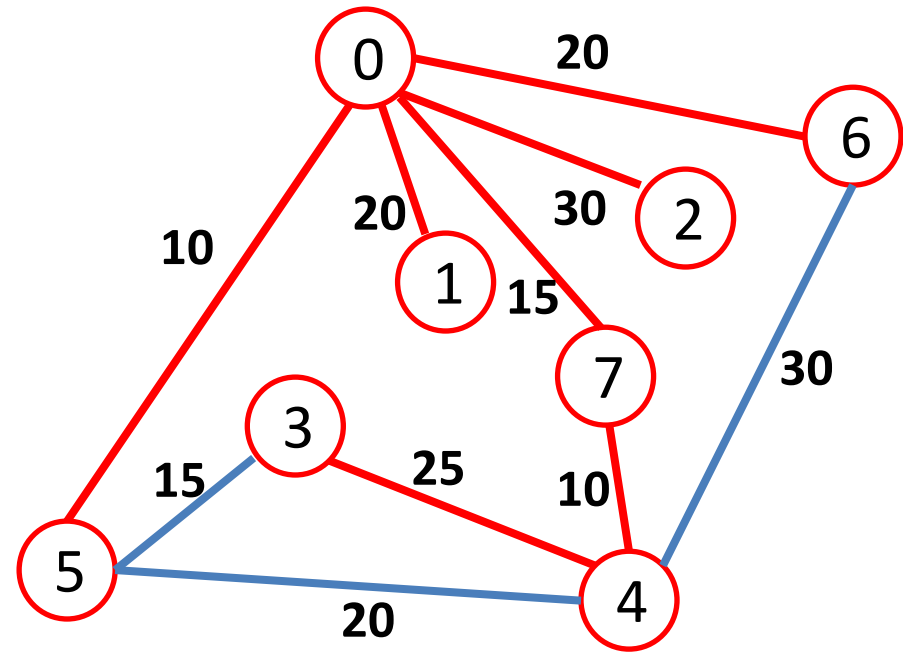
- Suppose we want to compute the SPST for vertex 7.
- Steps 7, 8, 9: $v = 6$
- Step 10: For $w =$ empty list



vertex	0	1	2	3	4	5	6	7
wt	15	35	45	35	10	25	35	0
st	7	0	0	4	7	0	0	7
in	1	1	0	1	1	1	1	1

Dijkstra Example

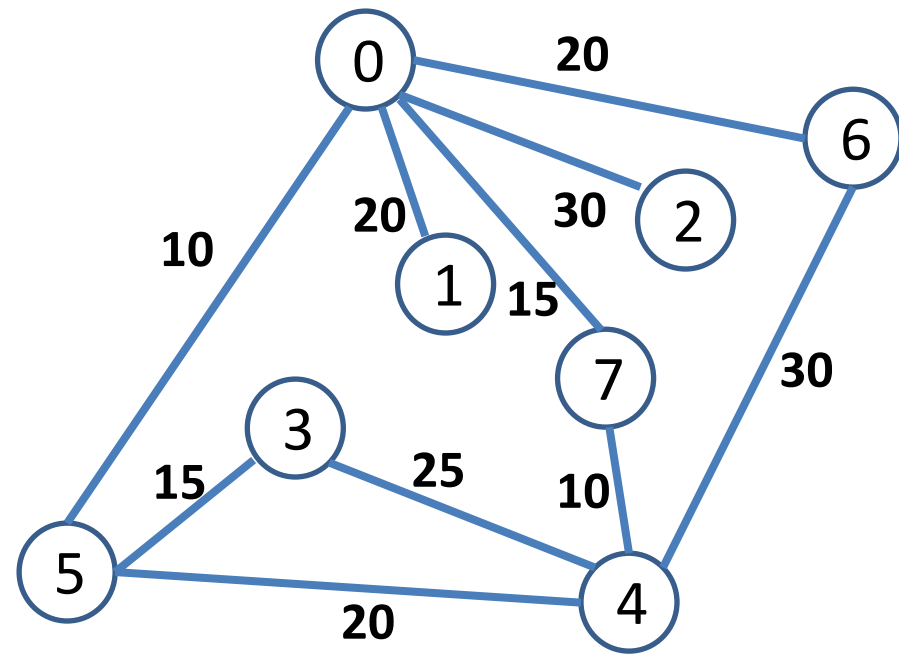
- Suppose we want to compute the SPST for vertex 7.
- Steps 7, 8, 9: $v = 6$
- Step 10: For $w =$ empty list



vertex	0	1	2	3	4	5	6	7
wt	15	35	45	35	10	25	35	0
st	7	0	0	4	7	0	0	7
in	1	1	1	1	1	1	1	1

Dijkstra Example

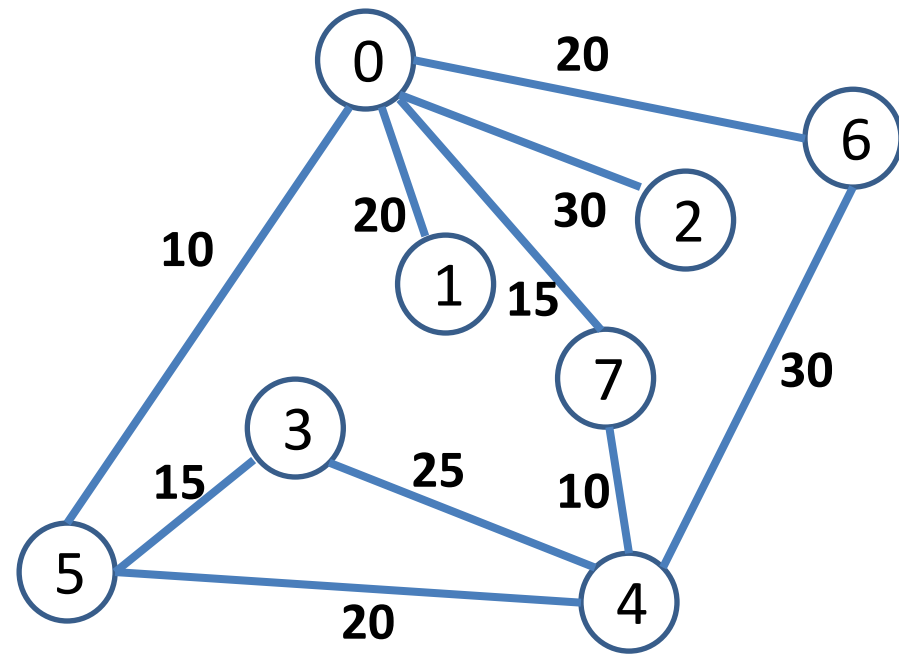
- Suppose we want to compute the SPST for vertex 4.
- First, we initialize arrays wt, st, in (steps 2, 3, 4).



vertex	0	1	2	3	4	5	6	7
wt	inf	inf	inf	inf	inf	inf	inf	inf
st	-1	-1	-1	-1	-1	-1	-1	-1
in	0	0	0	0	0	0	0	0

Dijkstra Example

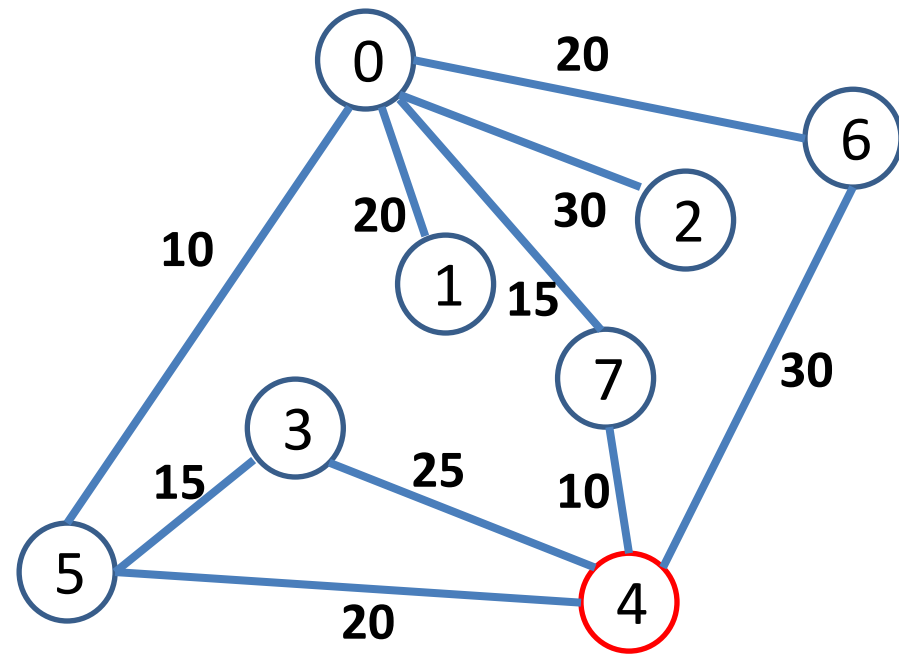
- Suppose we want to compute the SPST for vertex 4.
- Step 5.



vertex	0	1	2	3	4	5	6	7
wt	inf	inf	inf	inf	0	inf	inf	inf
st	-1	-1	-1	-1	4	-1	-1	-1
in	0	0	0	0	0	0	0	0

Dijkstra Example

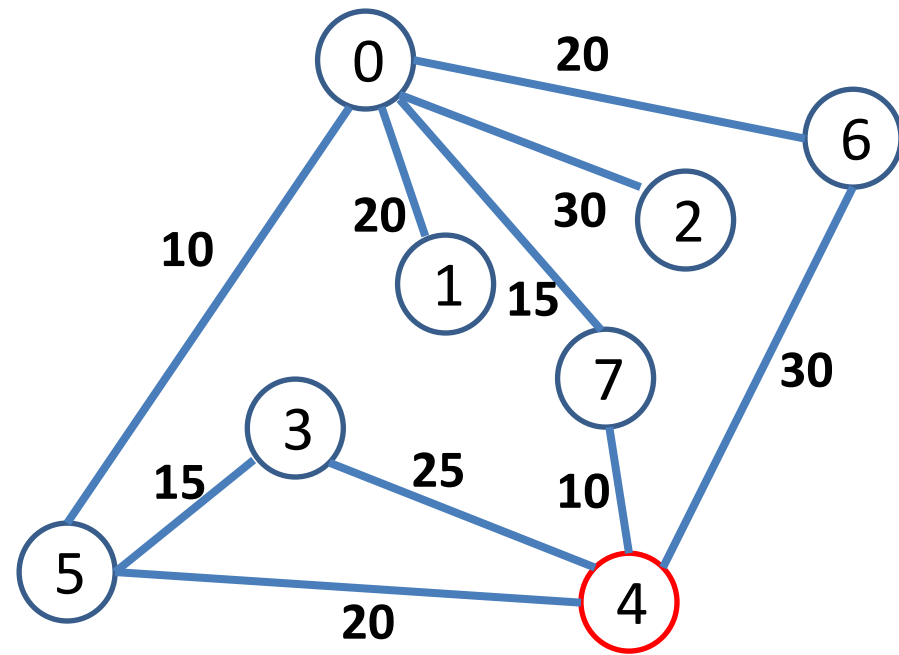
- Suppose we want to compute the SPST for vertex 4.
- Steps 7, 8, 9: $v = 4$



vertex	0	1	2	3	4	5	6	7
wt	inf	inf	inf	inf	0	inf	inf	inf
st	-1	-1	-1	-1	4	-1	-1	-1
in	0	0	0	0	1	0	0	0

Dijkstra Example

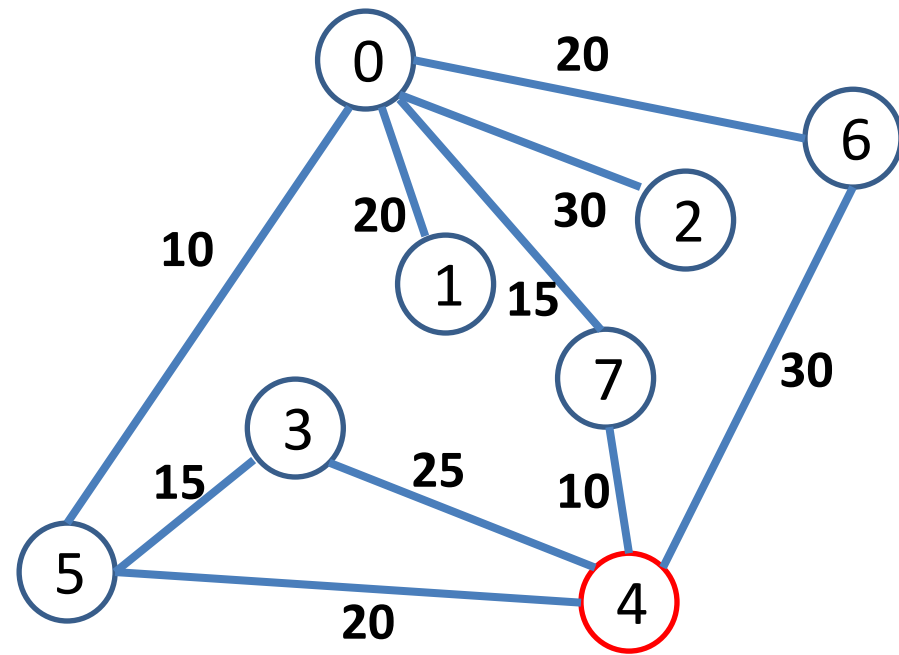
- Suppose we want to compute the SPST for vertex 4.
- Steps 7, 8, 9: $v = 4$
- Step 10: For $w = \{3, 5, 6, 7\}$
 - Step 11: Compare inf with 25
 - Steps 12, 13: $wt[3] = wt[4] + 25$, $st[3] = 4$.



vertex	0	1	2	3	4	5	6	7
wt	inf	inf	inf	25	0	inf	inf	inf
st	-1	-1	-1	4	4	-1	-1	-1
in	0	0	0	0	1	0	0	0

Dijkstra Example

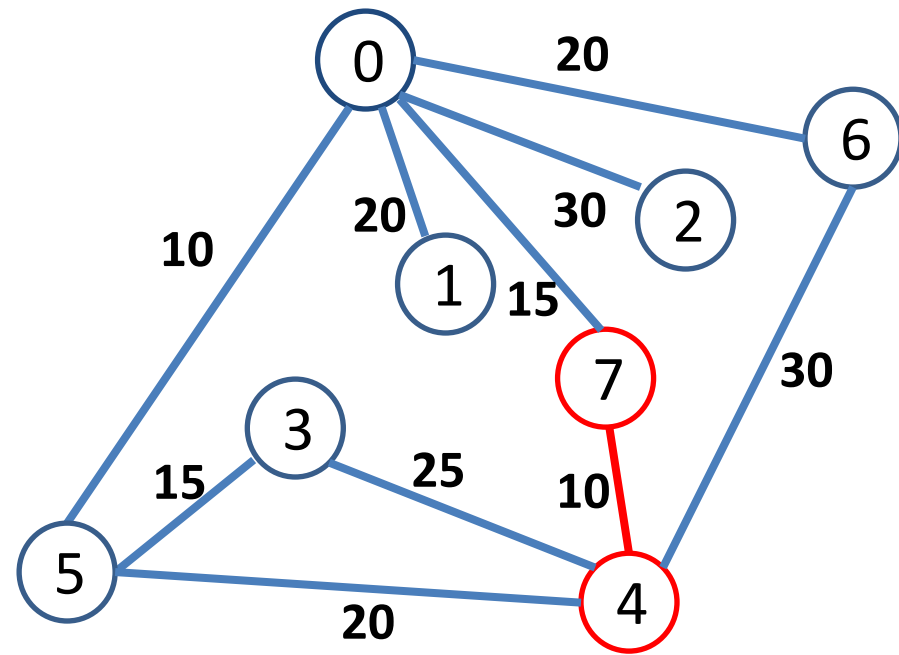
- Suppose we want to compute the SPST for vertex 4.
- Steps 7, 8, 9: $v = 4$
- Step 10: For $w = \{3, 5, 6, 7\}$
 - Steps 12, 13: update $wt[w]$, $st[w]$



vertex	0	1	2	3	4	5	6	7
wt	inf	inf	inf	25	0	20	30	10
st	-1	-1	-1	4	4	4	4	4
in	0	0	0	0	1	0	0	0

Dijkstra Example

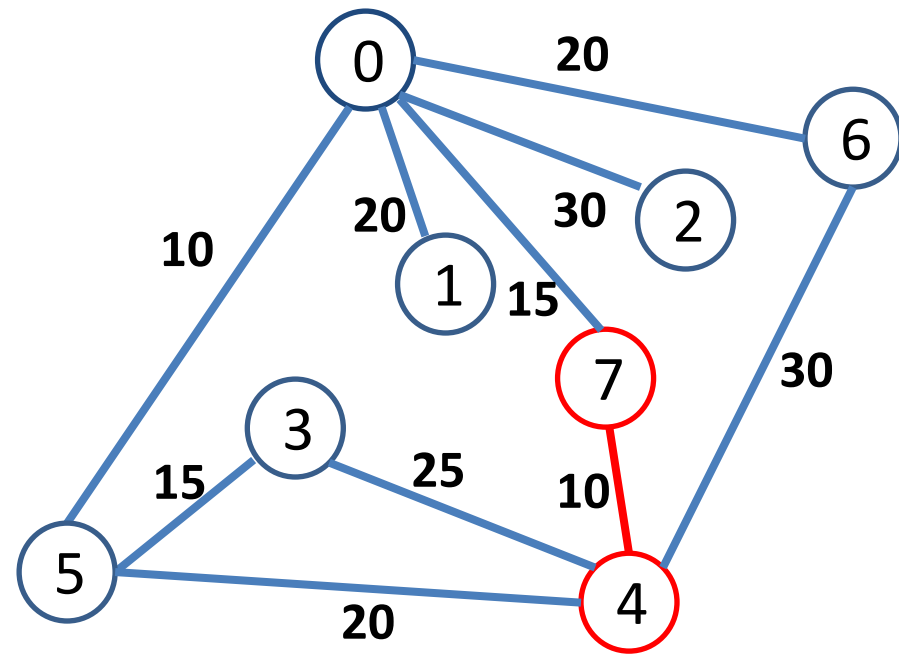
- Suppose we want to compute the SPST for vertex 4.
- Steps 7, 8, 9: $v = 7$



vertex	0	1	2	3	4	5	6	7
wt	inf	inf	inf	25	0	20	30	10
st	-1	-1	-1	4	4	4	4	4
in	0	0	0	0	1	0	0	1

Dijkstra Example

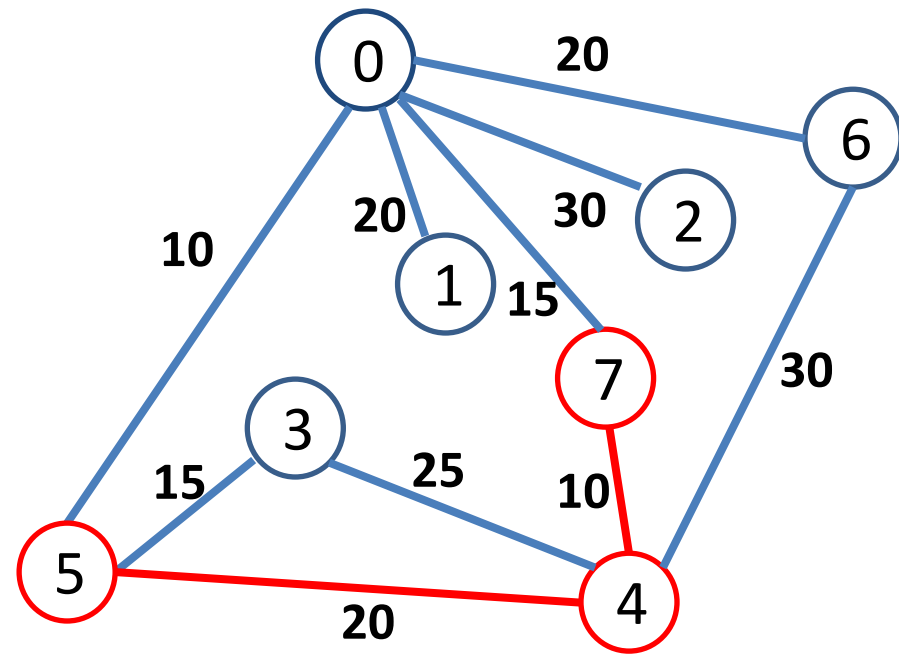
- Suppose we want to compute the SPST for vertex 4.
- Steps 7, 8, 9: $v = 7$
- Step 10: For $w = \{0\}$
 - Step 11: Compare inf with $10 + 15 = 25$.
 - Steps 12, 13: $\text{wt}[0] = \text{wt}[7] + 15$, $\text{st}[0] = 7$.



vertex	0	1	2	3	4	5	6	7
wt	25	inf	inf	25	0	20	30	10
st	7	-1	-1	4	4	4	4	4
in	0	0	0	0	1	0	0	1

Dijkstra Example

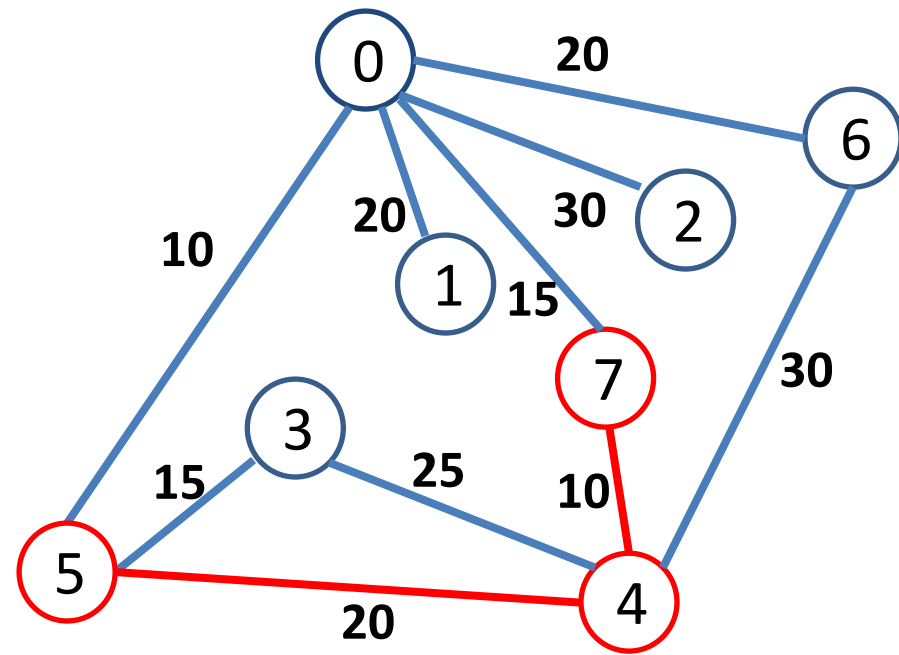
- Suppose we want to compute the SPST for vertex 4.
- Steps 7, 8, 9: $v = 5$



vertex	0	1	2	3	4	5	6	7
wt	25	inf	inf	25	0	20	30	10
st	7	-1	-1	4	4	4	4	4
in	0	0	0	0	1	1	0	1

Dijkstra Example

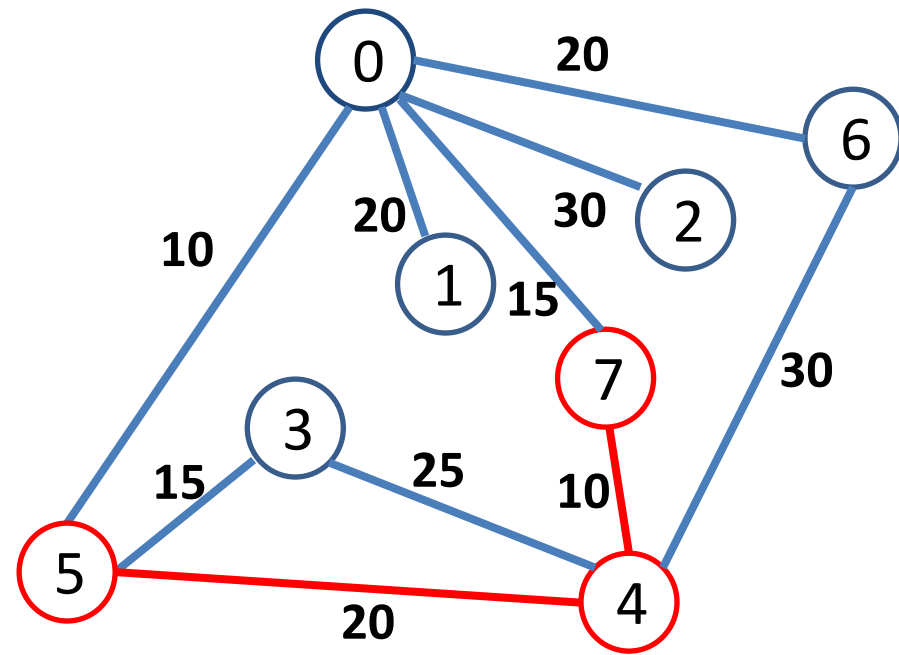
- Suppose we want to compute the SPST for vertex 4.
- Steps 7, 8, 9: $v = 5$
- Step 10: For $w = \{0, 3\}$
 - Step 11: Compare 25 with $20 + 10 = 25$.
NO UPDATE



vertex	0	1	2	3	4	5	6	7
wt	25	inf	inf	25	0	20	30	10
st	7	-1	-1	4	4	4	4	4
in	0	0	0	0	1	1	0	1

Dijkstra Example

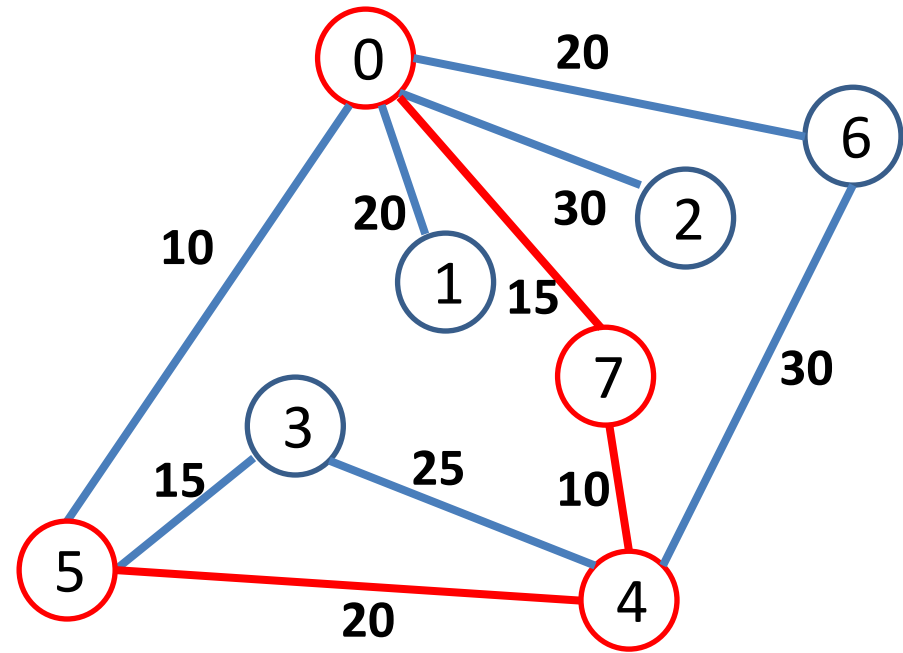
- Suppose we want to compute the SPST for vertex 4.
- Steps 7, 8, 9: $v = 5$
- Step 10: For $w = \{0, 3\}$
 - Step 11: Compare 25 with $20 + 15 = 35$.
NO UPDATE



vertex	0	1	2	3	4	5	6	7
wt	25	inf	inf	25	0	20	30	10
st	7	-1	-1	4	4	4	4	4
in	0	0	0	0	1	1	0	1

Dijkstra Example

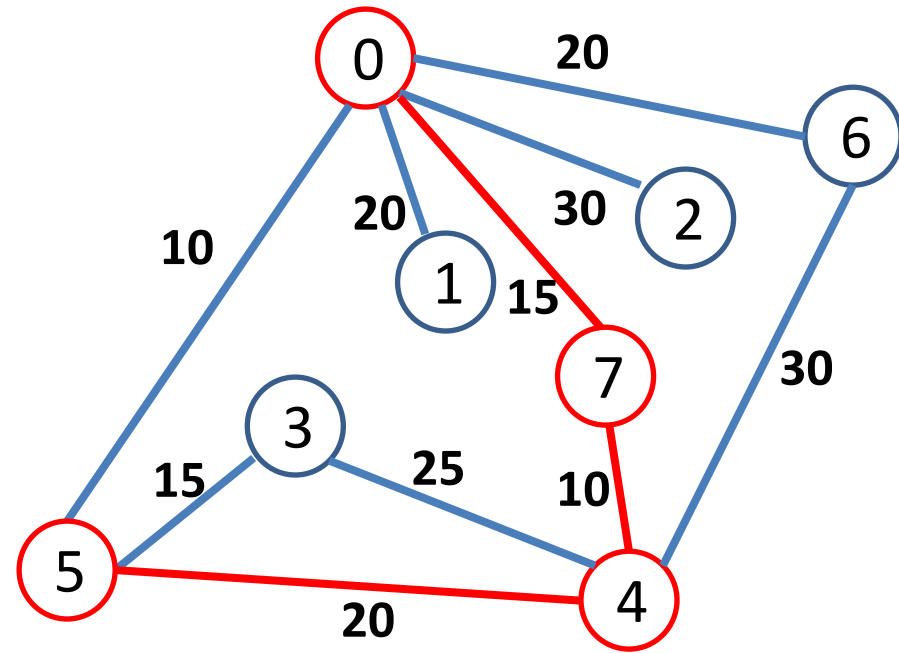
- Suppose we want to compute the SPST for vertex 4.
- Steps 7, 8, 9: $v = 0$



vertex	0	1	2	3	4	5	6	7
wt	25	inf	inf	25	0	20	30	10
st	7	-1	-1	4	4	4	4	4
in	1	0	0	0	1	1	0	1

Dijkstra Example

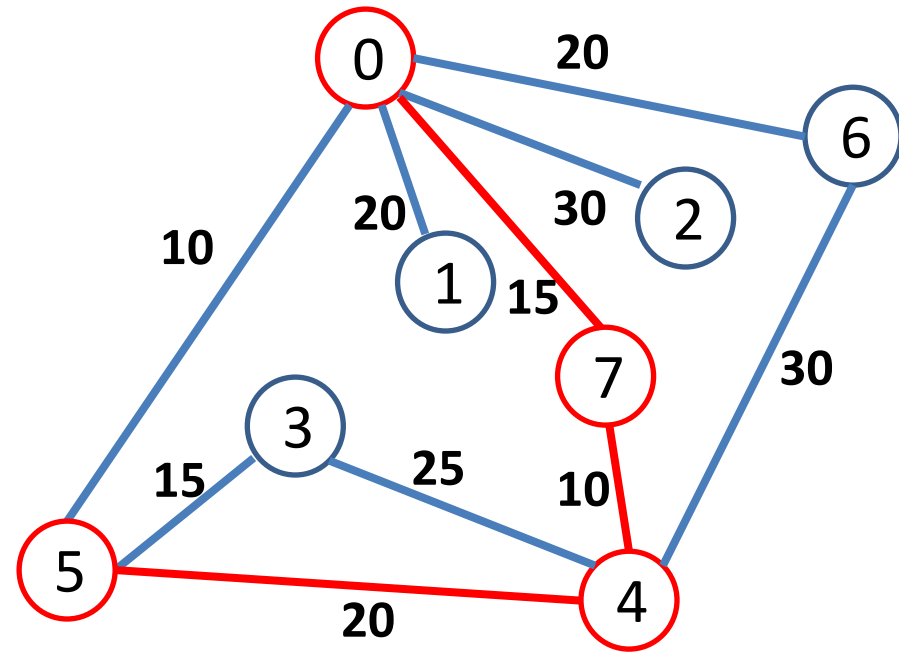
- Suppose we want to compute the SPST for vertex 4.
- Steps 7, 8, 9: $v = 0$
- Step 10: For $w = \{1, 2, 6\}$
 - Step 11: Compare inf with $25 + 20 = 45$.
 - Steps 12, 13: $wt[1] = wt[0] + 20$, $st[1] = 0$.



vertex	0	1	2	3	4	5	6	7
wt	25	45	inf	25	0	20	30	10
st	7	0	-1	4	4	4	4	4
in	1	0	0	0	1	1	0	1

Dijkstra Example

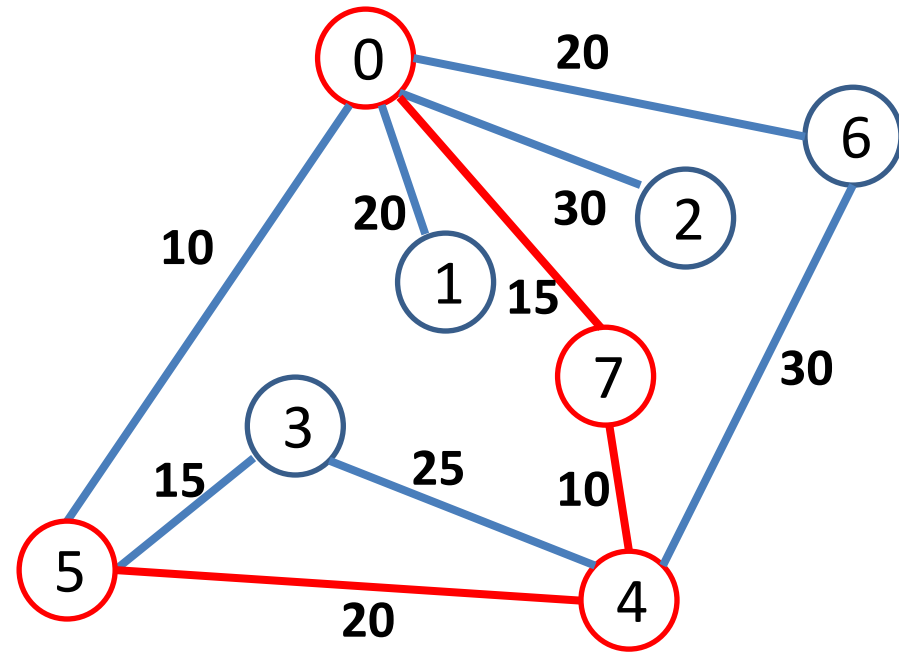
- Suppose we want to compute the SPST for vertex 4.
- Steps 7, 8, 9: $v = 0$
- Step 10: For $w = \{1, 2, 6\}$
 - Step 11: Compare inf with $25 + 30 = 55$.
 - Steps 12, 13: $wt[2] = wt[0] + 30$, $st[2] = 0$.



vertex	0	1	2	3	4	5	6	7
wt	25	45	55	25	0	20	30	10
st	7	0	0	4	4	4	4	4
in	1	0	0	0	1	1	0	1

Dijkstra Example

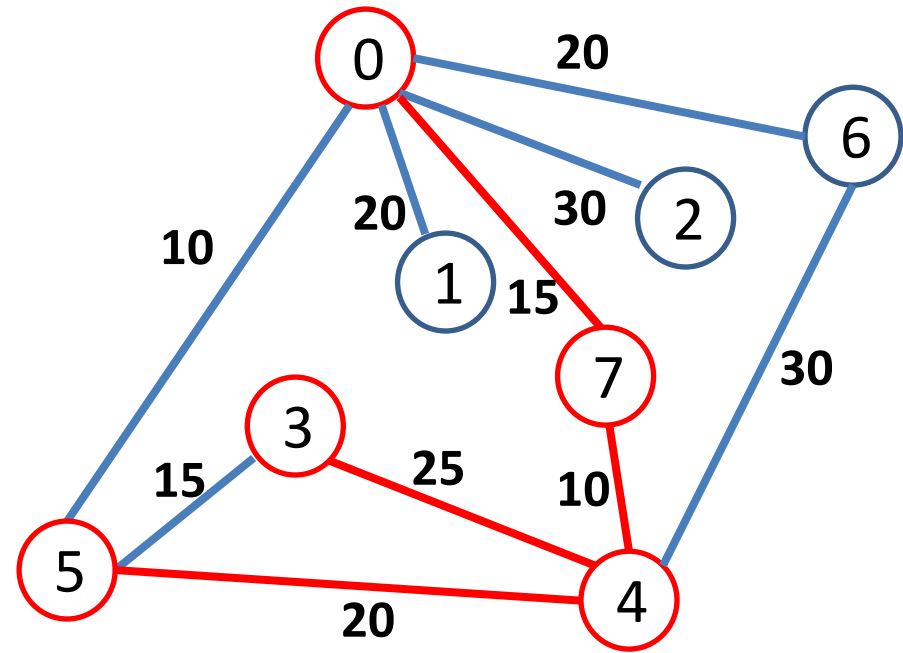
- Suppose we want to compute the SPST for vertex 4.
- Steps 7, 8, 9: $v = 0$
- Step 10: For $w = \{1, 2, 6\}$
 - Step 11: Compare 30 with $25 + 20 = 45$.
NO UPDATE



vertex	0	1	2	3	4	5	6	7
wt	25	45	55	25	0	20	30	10
st	7	0	0	4	4	4	4	4
in	1	0	0	0	1	1	0	1

Dijkstra Example

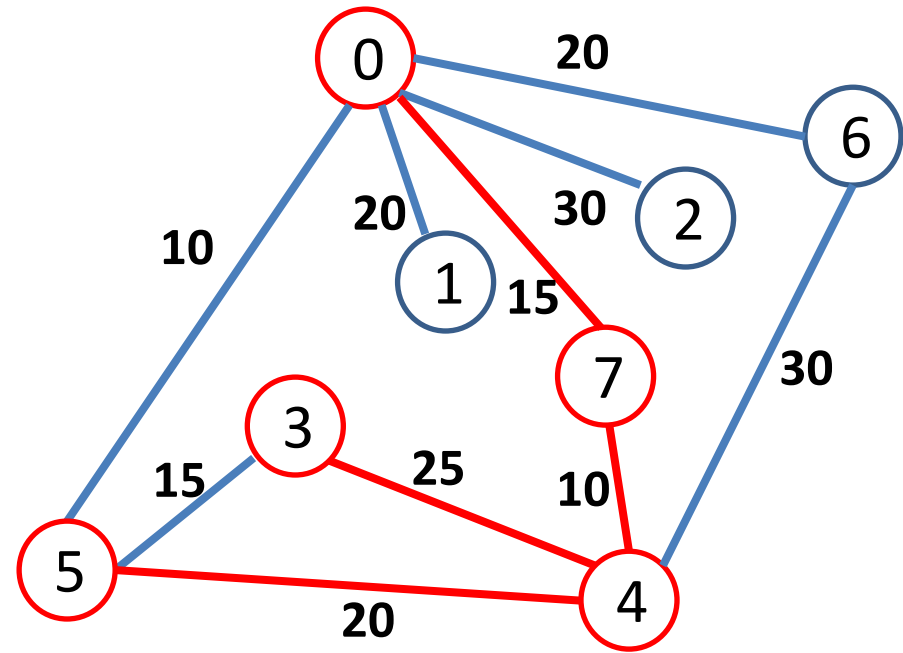
- Suppose we want to compute the SPST for vertex 4.
- Steps 7, 8, 9: $v = 3$



vertex	0	1	2	3	4	5	6	7
wt	25	45	55	25	0	20	30	10
st	7	0	0	4	4	4	4	4
in	1	0	0	1	1	1	0	1

Dijkstra Example

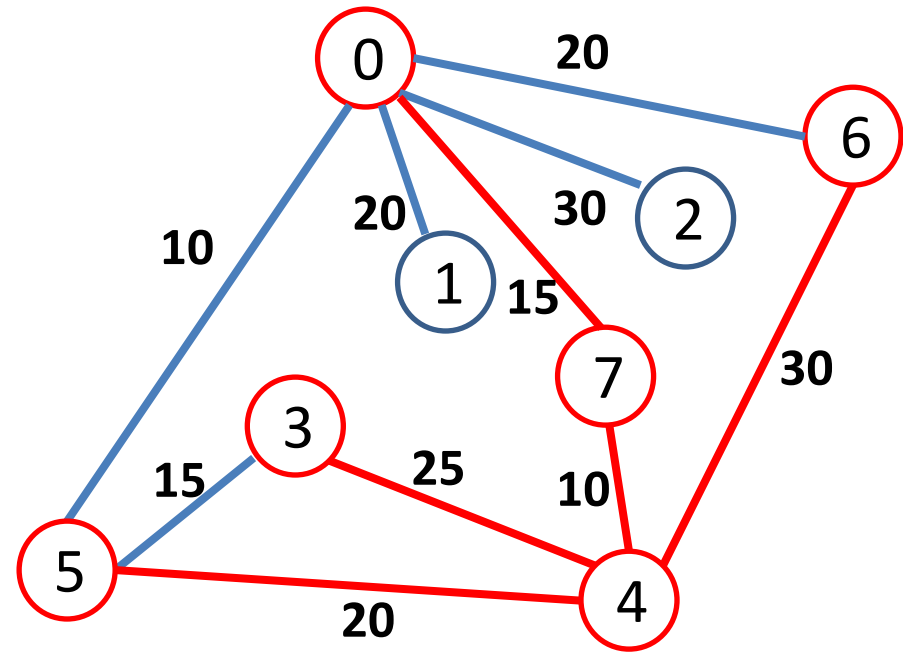
- Suppose we want to compute the SPST for vertex 4.
- Steps 7, 8, 9: $v = 3$
- Step 10: empty list



vertex	0	1	2	3	4	5	6	7
wt	25	45	55	25	0	20	30	10
st	7	0	0	4	4	4	4	4
in	1	0	0	1	1	1	0	1

Dijkstra Example

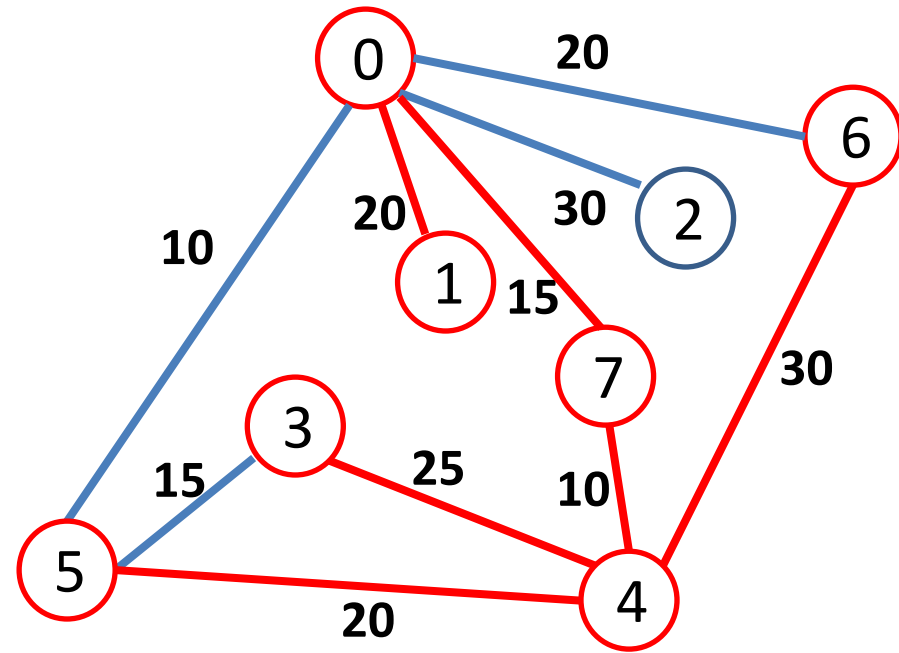
- Suppose we want to compute the SPST for vertex 4.
- Steps 7, 8, 9: $v = 6$
- Step 10: empty list



vertex	0	1	2	3	4	5	6	7
wt	25	45	55	25	0	20	30	10
st	7	0	0	4	4	4	4	4
in	1	0	0	1	1	1	1	1

Dijkstra Example

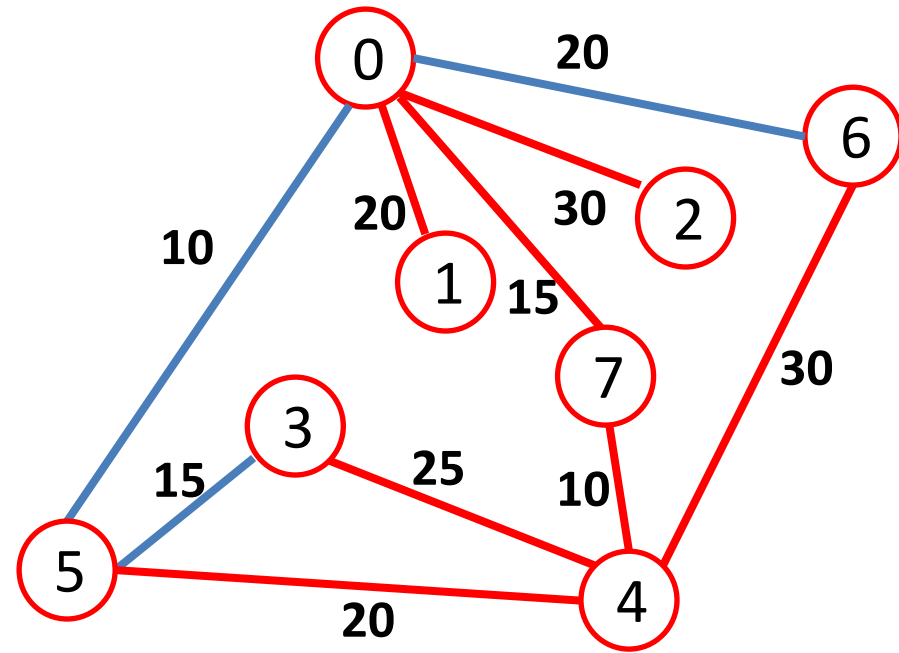
- Suppose we want to compute the SPST for vertex 4.
- Steps 7, 8, 9: $v = 1$
- Step 10: empty list



vertex	0	1	2	3	4	5	6	7
wt	25	45	55	25	0	20	30	10
st	7	0	0	4	4	4	4	4
in	1	1	0	1	1	1	1	1

Dijkstra Example

- Suppose we want to compute the SPST for vertex 4.
- Steps 7, 8, 9: $v = 2$
- Step 10: empty list



vertex	0	1	2	3	4	5	6	7
wt	25	45	55	25	0	20	30	10
st	7	0	0	4	4	4	4	4
in	1	1	1	1	1	1	1	1

All-Pairs Shortest Paths

- Before we describe an algorithm for computing the shortest paths among all pairs of vertices, we should agree on what this algorithm should return.
- We need to compute two $V \times V$ arrays:
 - $\text{dist}[v][w]$ is the distance of the shortest path from v to w .
 - $\text{path}[v][w]$ is the vertex following v , on the shortest path from v to w .
- Given these two arrays (after our algorithm has completed), how can we recover the shortest path between some v and w ?

All-Pairs Shortest Paths

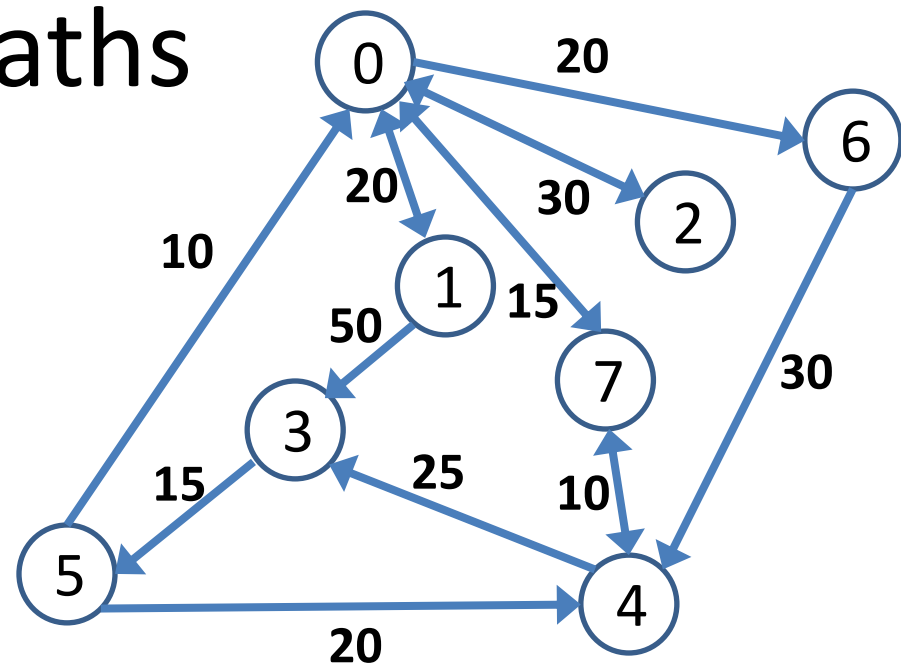
- We need to compute two $V \times V$ arrays:
 - $\text{dist}[v][w]$ is the distance of the shortest path from v to w .
 - $\text{path}[v][w]$ is the vertex following v , on the shortest path from v to w .
- Given these two arrays (after our algorithm has completed), how can we recover the shortest path between some v and w ?
- $\text{path} = \text{empty list}$
- $c = v$
- $\text{while}(\text{true})$
 - $\text{insert_to_end}(\text{path}, c)$
 - $\text{if } (c == w) \text{ break}$
 - $c = \text{path}[c][w]$

Computing Shortest Paths

- Overview: we can simply call Dijkstra's algorithm on each vertex.
- Time: V times the time of running Dijkstra's algorithm once.
 - $O(E \lg V)$ for one vertex.
 - $O(VE \lg V)$ for all vertices.
 - $O(V^3 \lg V)$ for dense graphs.
- There is a better algorithm for dense graphs, Floyd's algorithm, with $O(V^3)$ complexity, but we will not cover it.

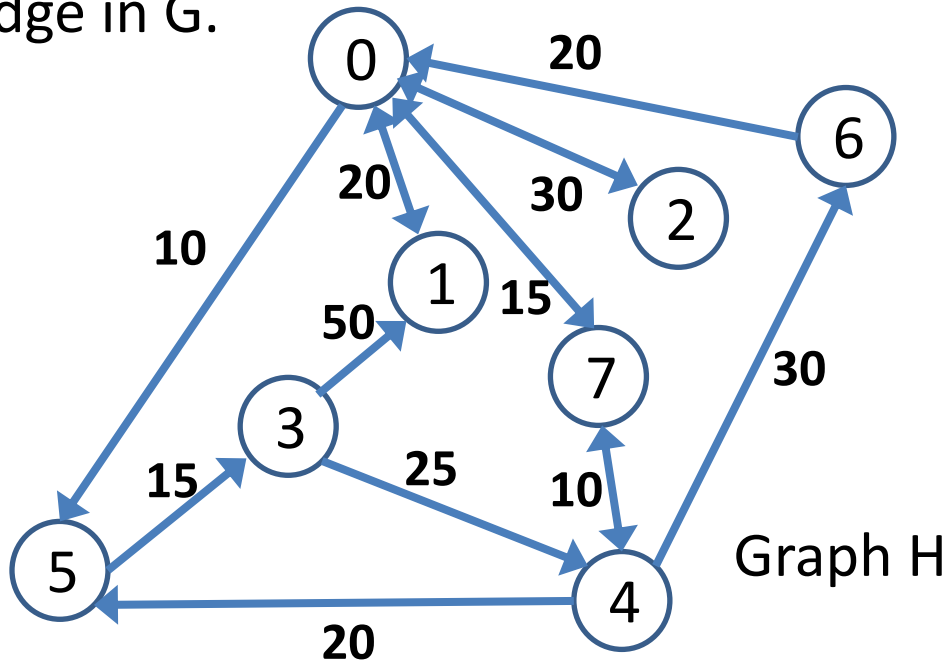
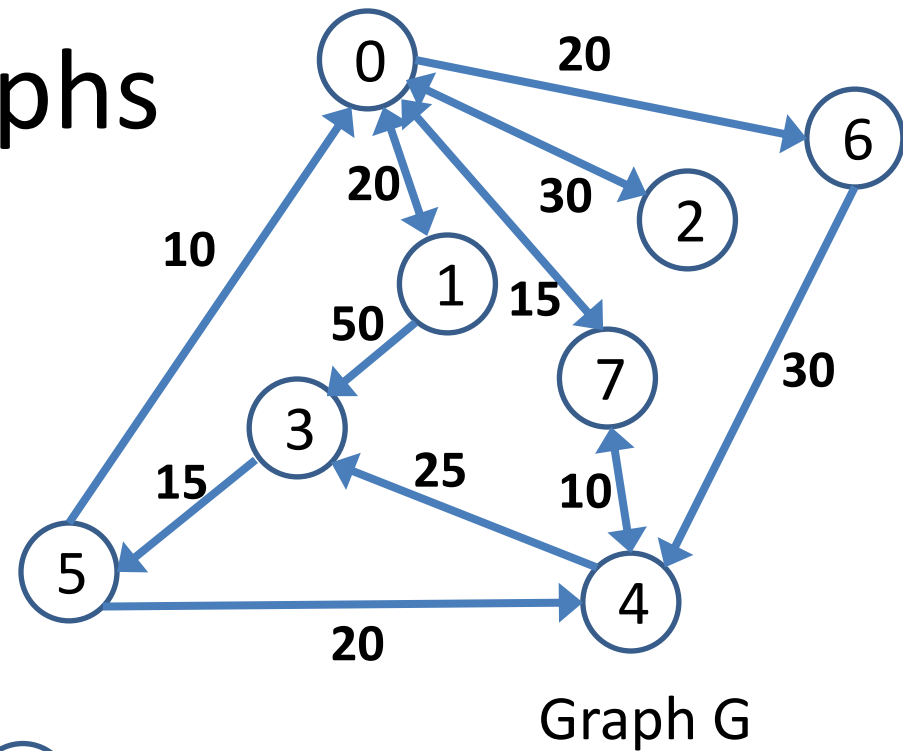
All-Pairs Shortest Paths Using Dijkstra

- The complete all-pairs algorithm is more complicated than simply calling Dijkstra's algorithm V times.
- Here is why:
- Suppose we call Dijkstra's algorithm on vertex 1.
- The algorithm computes arrays wt and st :
 - $wt[v]$: weight of shortest path from vertex 1 to v .
 - $st[v]$: parent vertex of v on shortest path from vertex 1 to v .
- How do arrays wt and st correspond to arrays $dist$ and $path$?
 - $dist[v][w]$ is the distance of the shortest path from v to w .
 - $path[v][w]$ is the vertex following v , on the shortest path from v to w .
- No useful correspondence!!!



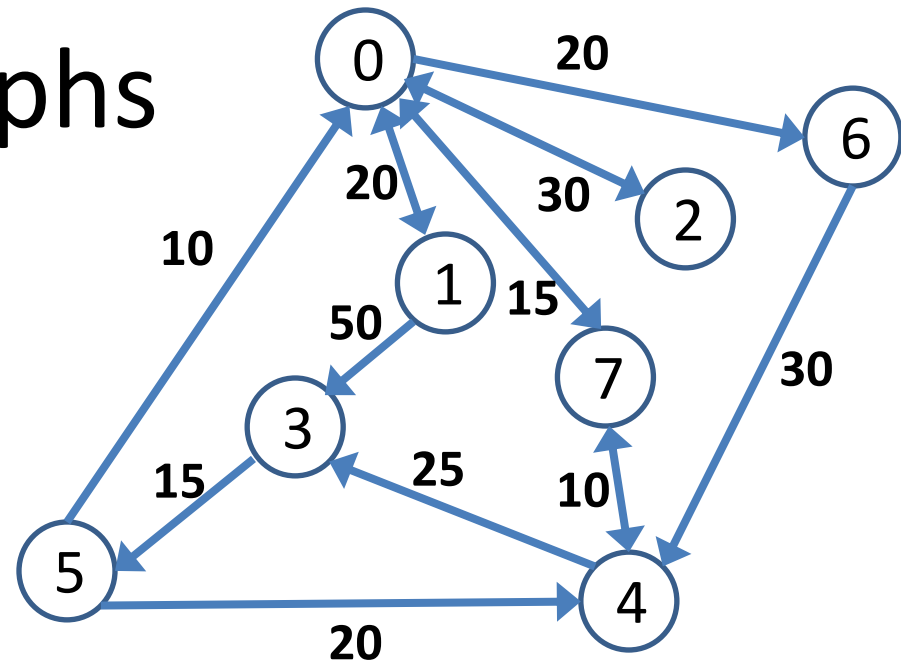
Using Reverse Graphs

- Suppose that G is the graph you see on the right.
- Suppose that H is the reverse graph, obtained by switching the direction of every single edge in G .



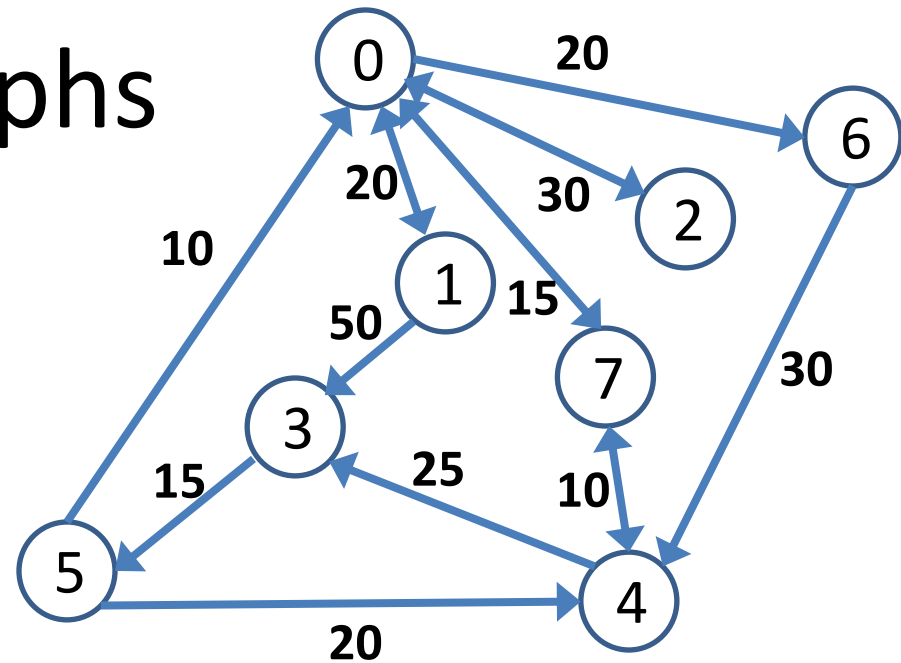
Using Reverse Graphs

- Suppose that G is the graph you see on the right.
- Suppose that H is the reverse graph, obtained by switching the direction of every single edge in G .
- Then, for any vertices v and w , the shortest path from w to v in H is simply the reverse of the shortest path from v to w in G .
- For example:
 - Shortest path from 1 to 4 in G :
 - Shortest path from 4 to 1 in H :



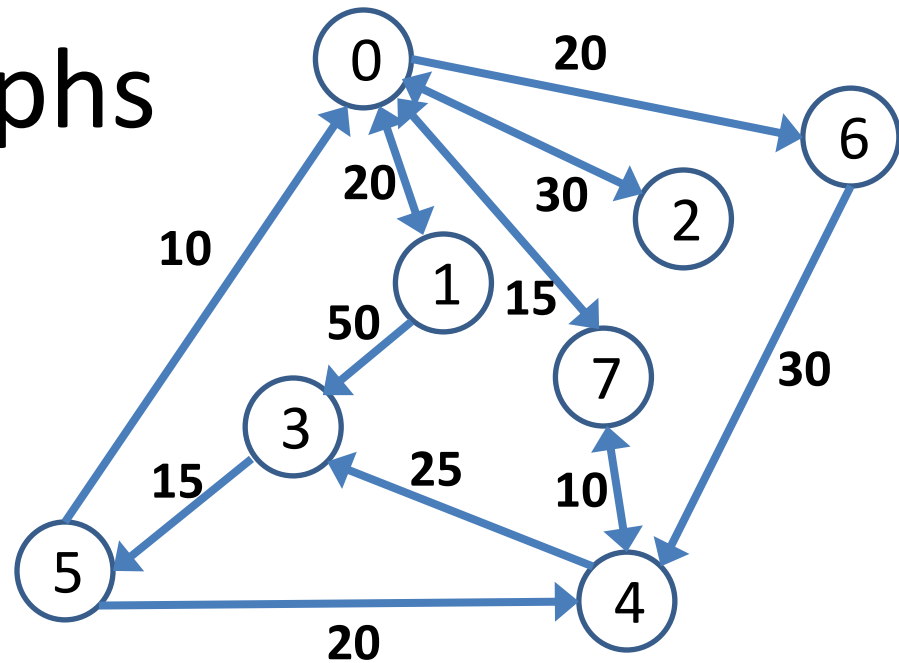
Using Reverse Graphs

- Suppose that G is the graph you see on the right.
- Suppose that H is the reverse graph, obtained by switching the direction of every single edge in G .
- Then, for any vertices v and w , the shortest path from w to v in H is simply the reverse of the shortest path from v to w in G .
- For example:
 - Shortest path from 1 to 4 in G : 1, 0, 7, 4
 - Shortest path from 4 to 1 in H : 4, 7, 0, 1.
 - These two paths are just reversed forms of each other, and they have the same weights.



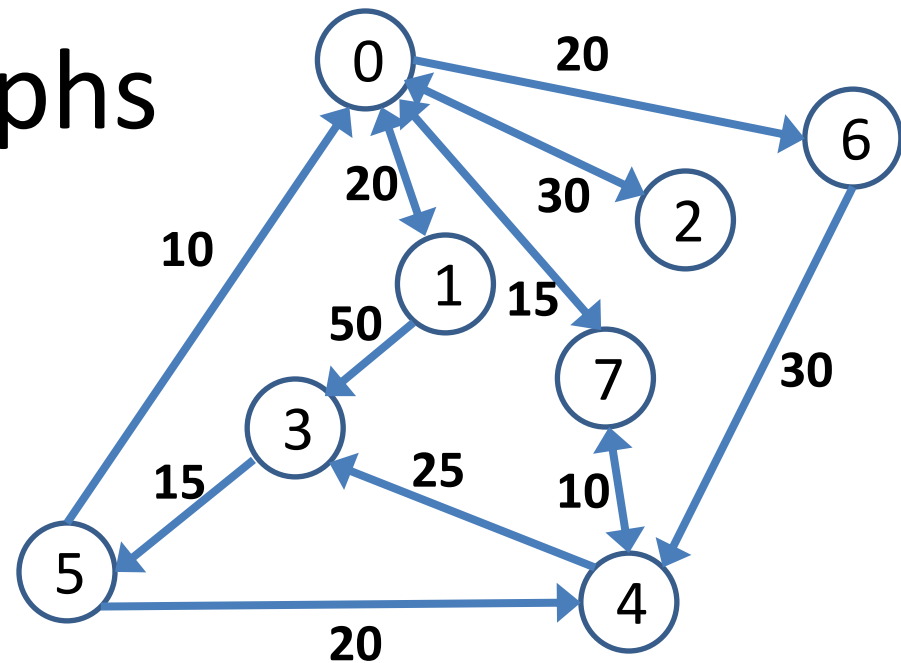
Using Reverse Graphs

- Suppose that we call Dijkstra's algorithm with source = vertex 1, on graph H (the **reverse graph** of what you see on the right).
- Consider the arrays wt and st we get as a result of that.
- These arrays are related to arrays dist and path on the **original graph G** (what you actually see on the right) as follows:
 - $\text{dist}[v][1] = \text{wt}[v]$.
 - $\text{path}[v][1] = \text{st}[v]$.
- Why?



Using Reverse Graphs

- Suppose that we call Dijkstra's algorithm with source = vertex 1, on graph H (the reverse graph of what you see on the right).
- Consider the arrays wt and st we get as a result of that.
- wt[v] is the weight of the shortest path from 1 to v in H.
 - Therefore, wt[v] is the weight of the shortest path from v to 1 in G.
 - Therefore, dist[v][1] = wt[v].
- st[v] is the parent of v on the shortest path from 1 to v in H.
 - Therefore, st[v] is the vertex following v on the shortest path from v to 1 in G.
 - Therefore, path[v][1] = st[v].



Using Dijkstra's Algorithm for All-Pairs Shortest Paths

Input: graph G .

1. Construct reverse graph H .
2. For each s in $\{0, \dots, V-1\}$:
 3. Call Dijkstra's algorithm on graph H , with source = s .
 4. For each v in $\{0, \dots, V-1\}$:
 5. $\text{dist}[v][s] = \text{wt}[v]$.
 6. $\text{path}[v][s] = \text{st}[v]$.