

# CSE 3302


## Programming Languages

### Control I

#### Expressions and Statements


Chengkai Li, Weimin He

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## Control



- Control:
  - what gets executed, when, and in what order.
- Abstraction of control:
  - Expression
  - Statement
  - Exception Handling
  - Procedures and functions

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## Expression vs. Statement



- In pure (mathematical) form:
  - Expression:
    - no side effect
    - return a value
  - Statement:
    - side effect
    - no return value
- Functional languages aim at achieving this pure form
- No clear-cut in most languages

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## Expression



- Constructed recursively:
  - Basic expression (literal, identifiers)
  - Operators, functions, special symbols
- Number of operands:
  - unary, binary, ternary operators
- Operator, function: equivalent concepts
  - $(3+4)*5$  (infix notation)
  - $\text{mul}(\text{add}(3,4), 5)$ 
    - $^{**} (^{+}(3,4), 5)$  (Ada, prefix notation)
    - $(^* (+ 3 4) 5)$  (LISP, prefix notation)

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## Postfix notation




- PostScript:
  - %!PS
  - /Courier findfont
  - 20 scalefont
  - setfont
  - 72 500 moveto
  - (Hello world!) show
  - showpage

<http://en.wikipedia.org/wiki/PostScript>

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## Expression and Side Effects



- Side Effects:
  - changes to memory, input/output
  - Side effects can be undesirable
  - But a program without side effects does nothing!
- Expression:
  - No side effect: Order of evaluating subexpressions doesn't matter (mathematical forms)
  - Side effect: Order matters

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## Applicative Order Evaluation (Strict Evaluation)




- Evaluate the operands first, then apply operators (bottom-up evaluation) (subexpressions evaluated, no matter whether they are needed)

- But is 3+4 or 5-6 evaluated first?

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## Order Matters



**C:**

```
int x=1;
int f(void) {
  x=x+1;
  return x;
}
main(){
  printf("%d\n", x + f());
  return 0;
}
```

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**Java:**

```
class example
{ static int x = 1;
  public static int f()
  {
    x = x+1;
    return x;
  }
  public static void main(String[] args)
  {
    System.out.println(x+f());
  }
}
```


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Many languages don't specify the order, including C, java.

- C: usually right-to-left
- Java: always left-to-right, but not suggested to rely on that.

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## Expected Side Effect



- Assignment (**expression, not statement**)  
 $x = (y = z)$  (right-associative operator)


Why?

- x++, ++x**

```
int x=1;
int f(void) {
  return x++;
}
main(){
  printf("%d\n", x + f());
  return 0;
}
```

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## Sequence Operator




- (**expr1, expr2, ..., exprn**)
  - Left to right (this is indeed specified in C)
  - The return value is **exprn**

```
x=1;
y=2;
x = (x=x+1, y++, x+y);
printf("%d\n", x);
```

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## Non-strict evaluation



- Evaluating an expression without necessarily evaluating all the subexpressions.
- short-circuit Boolean expression
- if-expression, case-expression

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## Short-Circuit Evaluation



- if (false and x) ... if (true or x) ...
  - No need to evaluate x, no matter x is true or false
- What is it good for?
  - if (i <= lastindex and a[i] >= x) ...
  - if (p != NULL and p->next == q) ...
- Ada: allow both short-circuit and non short-circuit.
  - if (x /= 0) and then (y/x > 2) then ...
  - if (x /= 0) and (y/x > 2) then ... ?
  - if (ptr = null) or else (ptr.x = 0) then ... ?
  - if (ptr = null) or (ptr.x = 0) then ... ?

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## if-expression



- `if (test-exp, then-exp, else-exp)`  
ternary operator
  - test-exp is always evaluated first
  - Either then-exp or else-exp are evaluated, not both
- `if e1 then e2 else e3` (ML)
- `e1 ? e2 : e3` (C)
- Different from if-statement?

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## case-expression



- ML:  
case color of  
  red => "R" |  
  blue => "B" |  
  green => "G" |  
  \_ => "AnyColor";

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## Normal order evaluation (lazy evaluation)



- When there is no side-effect:  
Normal order evaluation (Expressions evaluated in mathematical form)
  - Operation evaluated *before* the operands are evaluated;
  - Operands *evaluated only when necessary*.
- `int double (int x) { return x+x; }`  
`int square (int x) { return x*x; }`
- Applicative order evaluation : `square(double(2)) = ...`  
Normal order evaluation : `square(double(2)) = ...`

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## What is it good for?



```
(p!=NULL) ? p->next : NULL
      ↓
int if_exp(bool x, int y, int z)
{ if (x)
  return y;
  else
  return z;
}
if_exp(p!=NULL, p->next, NULL);
```

- With side effect, it may hurt you:  
`int get_int(void) {`  
  `int x;`  
  `scanf("%d" &x);`  
  `return x;`  
}

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## Examples



- Call by Name (Algol60)
  - Macro
- ```
#define swap(a, b) {int t; t = a; a = b; b = t;}
```

– What are the problems here?

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## Unhygienic Macros



- Call by Name (Algol60)
  - Macro
- ```
#define swap(a, b) {int t; t = a; a = b; b = t;}

main () {
  int t=2;
  int s=5;
  swap(s,t);
}

main () {
  int t=2;
  int s=5;
  {int t; t = s; s = t; t = t;}
}
```
- 
- ```
#define DOUBLE(x) {x+x;}

main () {
  int a;
  a = DOUBLE(get_int());
  printf("a=%d\n", a);
}

main () {
  int a;
  a = get_int()+get_int();
  printf("a=%d\n", a);
}
```

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## Statements



- If-statements, case-(switch-)statements, loops