


CSE 3302 Programming Languages




Logic Programming: Prolog

Chengkai Li
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SWI-Prolog



- <http://www.swi-prolog.org/>
- Available for:
Linux, Windows, MacOS

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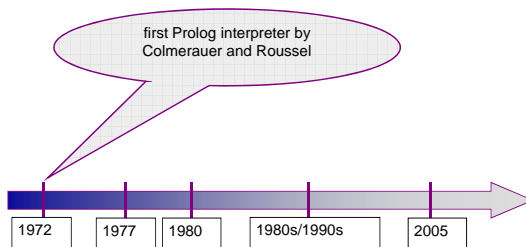
Prolog



- Prolog:
“Programming in Logic” (PROgrammation en LOGique)
- One (and maybe the only one) successful logic programming languages
- Useful in AI applications, expert systems, natural language processing, database query languages
- Declarative instead of procedural: “What” instead of “How”


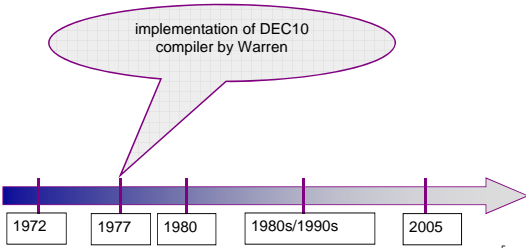
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History of Prolog


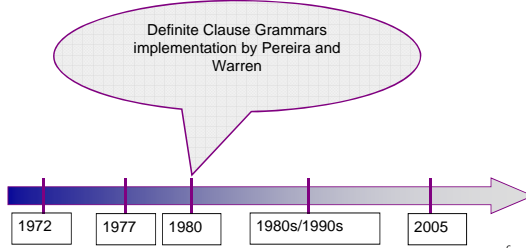
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History of Prolog

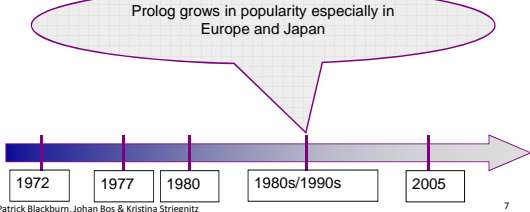
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History of Prolog

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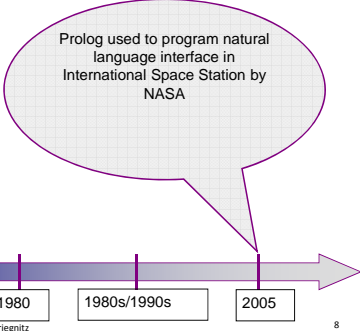
History of Prolog



Prolog grows in popularity especially in Europe and Japan

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History of Prolog



Prolog used to program natural language interface in International Space Station by NASA

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Logic Programming

- Program
Axioms (facts): true statements
- Input to Program
query (goal): statement true (theorems) or false?
- Thus
Logic programming systems = deductive databases
datalog

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Example

- Axioms:
0 is a natural number. (Facts)
For all x, if x is a natural number, then so is the successor of x.
- Query (goal).
Is 2 natural number? (can be proved by facts)
Is -1 a natural number? (cannot be proved)

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Another example

- Axioms:
The factorial of 0 is 1. (Facts)
If m is the factorial of n - 1, then n * m is the factorial of n.
- Query:
The factorial of 2 is 3?

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First-Order Predicate Calculus

- Logic used in logic programming:
First-order predicate calculus
First-order predicate logic
Predicate logic
First-order logic

$$\forall x (x \neq x+1)$$

- Second-order logic
 $\forall S \forall x (x \in S \vee x \notin S)$

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First-Order Predicate Calculus: Example

- natural(0)
 $\forall X, \text{natural}(X) \rightarrow \text{natural}(\text{successor}(x))$
- $\forall X$ and $Y, \text{parent}(X,Y) \rightarrow \text{ancestor}(X,Y)$.
 $\forall A, B,$ and $C, \text{ancestor}(A,B)$ and $\text{ancestor}(B,C) \rightarrow \text{ancestor}(A,C)$.
 $\forall X$ and $Y, \text{mother}(X,Y) \rightarrow \text{parent}(X,Y)$.
 $\forall X$ and $Y, \text{father}(X,Y) \rightarrow \text{parent}(X,Y)$.
 $\text{father}(\text{bill}, \text{jill})$.
 $\text{mother}(\text{jill}, \text{sam})$.
 $\text{father}(\text{bob}, \text{sam})$.
- factorial(0,1).
 $\forall N$ and $M, \text{factorial}(N-1,M) \rightarrow \text{factorial}(N,N*M)$.

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First-Order Predicate Calculus: statements

Symbols in statements:

- Constants (a.k.a. atoms)**
 numbers (e.g., 0) or names (e.g., bill).
- Predicates**
 Boolean functions (true/false). Can have arguments. (e.g. $\text{parent}(X, Y)$).
- Functions**
 non-Boolean functions ($\text{successor}(X)$).
- Variables**
 e.g., X .
- Connectives (operations)**
 and, or, not
 implication (\rightarrow): $a \rightarrow b$ (**b or not a**)
 equivalence (\leftrightarrow): $a \leftrightarrow b$ ($a \rightarrow b$ and $b \rightarrow a$)

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First-Order Predicate Calculus: statements (cont'd)

- Quantifiers**
 universal quantifier "for all" \forall
 existential quantifier "there exists" \exists
 bound variable (a variable introduced by a quantifier)
 free variable
- Punctuation symbols**
 parentheses (for changing associativity and precedence.)
 comma
 period
- Arguments to predicates and functions can only be terms:**
 - Contain constants, variables, and functions.
 - Cannot have predicates, qualifiers, or connectives.

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Problem Solving

- Program = Data + Algorithms
- Program = Object.Message(Object)
- Program = Functions Functions
- Algorithm = Logic + Control

Programmers:
 facts/axioms/statements

Logic programming systems:
 prove goals from axioms

- The holy grail: we specify the logic itself, the system proves.
 - Not totally realized by logic programming languages. Programmers must be aware of how the system proves, in order to write efficient, or even correct programs.
- Prove goals from facts:
 - Resolution and Unification**

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Horn Clause

- First-order logic too complicated for an effective logic programming system.
- Horn Clause: a fragment of first-order logic
 $b \leftarrow a_1 \text{ and } a_2 \text{ and } a_3 \dots \text{ and } a_n$.

head

←

body

←

no "or" and no quantifier

$b \leftarrow \text{fact}$
 $\leftarrow b. \text{query}$

- Variables in head: universally quantified
 Variables in body only: existentially quantified
- Need "or" in head? Multiple clauses

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Horn Clauses: Example

- First-Order Logic:
 $\text{natural}(0)$.
 $\forall X, \text{natural}(X) \rightarrow \text{natural}(\text{successor}(x))$.

- Horn Clause:
 $\text{natural}(0)$.
 $\text{natural}(\text{successor}(x)) \leftarrow \text{natural}(X)$.

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Horn Clauses: Example



- First-Order Logic:

```
factorial(0,1).
 $\forall N \text{ and } \forall M, \text{ factorial}(N-1,M) \rightarrow \text{factorial}(N,N*M).$ 
```



- Horn Clause:

```
factorial(0,1).
factorial(N,N*M) ← factorial(N-1,M).
```

Horn Clauses: Example



- Horn Clause:

```
ancestor(X,Y) ← parent(X,Y).
ancestor(A,C) ← ancestor(A,B) and ancestor(B,C).
parent(X,Y) ← mother(X,Y).
parent(X,Y) ← father(X,Y).
father(bill,jill).
mother(jill,sam).
father(bob,sam).
```

Horn Clauses: Example



- First-Order Logic:

```
 $\forall X, \text{ mammal}(X) \rightarrow \text{legs}(X,2) \text{ or } \text{legs}(X,4).$ 
```



- Horn Clause:

```
legs(X,4) ← mammal(X) and not legs(X,2).
legs(X,2) ← mammal(X) and not legs(X,4).
```

Prolog syntax



- :- for ←
, for and

```
ancestor(X,Y) :- parent(X,Y).
ancestor(X,Y) :- ancestor(X,Z), ancestor(Z,Y).
parent(X,Y) :- mother(X,Y).
parent(X,Y) :- father(X,Y).
father(bill,jill).
mother(jill,sam).
father(bob,sam).
```

Resolution and Unification



Resolution



- Resolution: Using a clause, replace its head in the second clause by its body, if they “match”.

```
a ← a1, ..., an.
b ← b1, ..., bi, ..., bm.
```

if b_i matches a_j

```
b ← b1, ..., a1, ..., an, ..., bm.
```

Resolution: Another view

- Resolution: Combine two clauses, and cancel matching statements on both sides.
- $a \leftarrow a_1, \dots, a_n.$
- $b \leftarrow b_1, \dots, b_i, \dots, b_m.$

~~a~~ , $b \leftarrow a_1, \dots, a_n, b_1, \dots, \overline{b_i}, \dots, b_m.$

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Problem solving in logic programming systems

- Program:
 - Statements/Facts (clauses).
- Goals:
 - Headless clauses, with a list of **subgoals**.
- **Problem solving by resolution:**
 - Matching subgoals with the heads in the facts, and replacing the subgoals by the corresponding bodies.
 - Cancelling matching statements.
 - Recursively do this, till we eliminate all goals. (Thus original goals proved.)

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Example

- Program:


```
mammal(human).
```
- Goal:


```
← mammal(human).
```
- Proving:


```
mammal(human) ← mammal(human).
      ←.
```

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Example

- Program:


```
legs(X,2) ← mammal(X), arms(X,2).
legs(X,4) ← mammal(X), arms(X,0).
mammal(horse).
arms(horse,0).
```
- Goal:


```
← legs(horse,4).
```
- Proving: ?

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Unification

- Unification: Pattern matching to make statements identical (when there are variables).
- Set variables equal to patterns: **instantiated**.
- In previous example:
 - legs(X,4) and legs(horse,4) are unified.**
 - (X is instantiated with horse.)**

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Unification: Example

- Euclid's algorithm for greatest common divisor
- Program:


```
gcd(U,0,U).
gcd(U,V,W) ← not zero(V), gcd(V, U mod V, W).
```
- Goals:


```
← gcd(15,10,X).
```

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Things unspecified



- The order to resolve subgoals.
- The order to use clauses to resolve subgoals.
- Possible to implement systems that don't depend on the order, but too inefficient.
- Thus programmers must know the orders used by the language implementations. (Search Strategies)

Example



- Program:


```
ancestor(X,Y) :- ancestor(X,Z), parent(Z,Y).
ancestor(X,Y) :- parent(X,Y).
parent(X,Y) :- mother(X,Y).
parent(X,Y) :- father(X,Y).
father(bill,jill).
mother(jill,sam).
father(bob,sam).
```
- Goals:


```
← ancestor(bill,sam).
← ancestor(X,bob).
```